

An Algorithmic Meta Theorem for Homomorphism Indistinguishability

AlMoTh 2024, 29 February 2024

Tim Seppelt



Research Training Group –
Uncertainty and Randomness
in Algorithms, Verification,
and Logic

RWTHAACHEN
UNIVERSITY



Chair for Logic
and Theory of
Discrete Systems

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Complexity of Homomorphism Indistinguishability

$\text{HOMIND}(\mathcal{F})$

Input Graphs G and H .

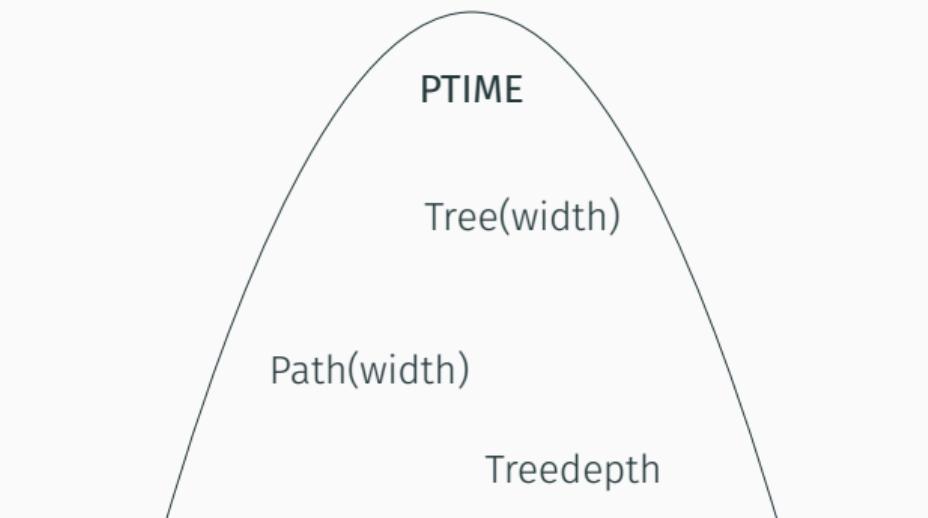
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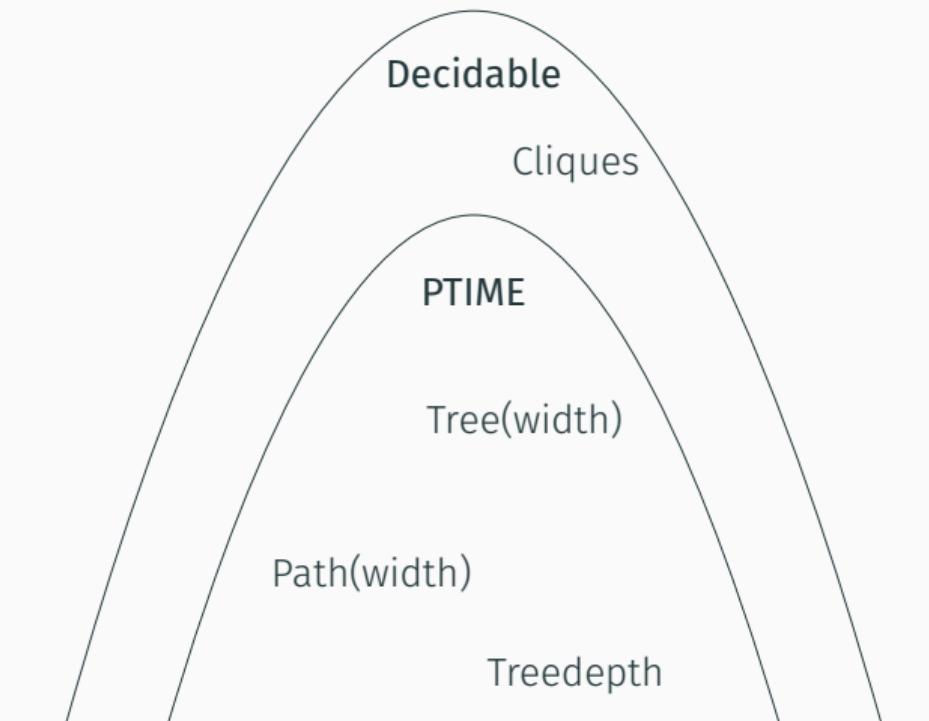
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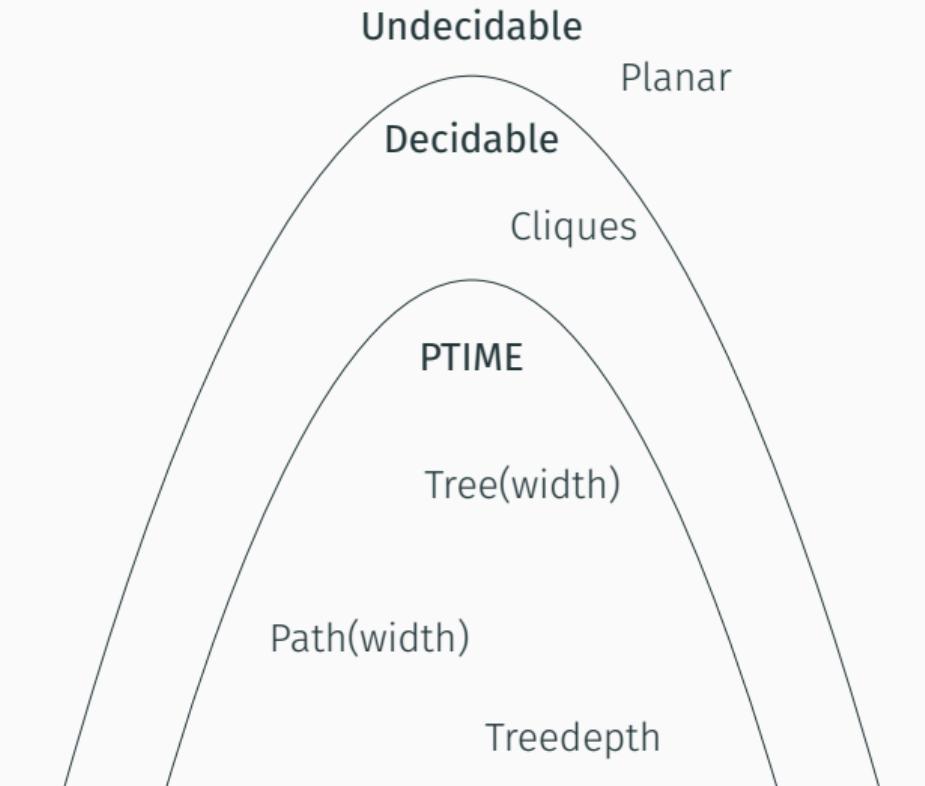
Complexity of Homomorphism Indistinguishability

Mančinska and Roberson (2020); Böker et al. (2019)

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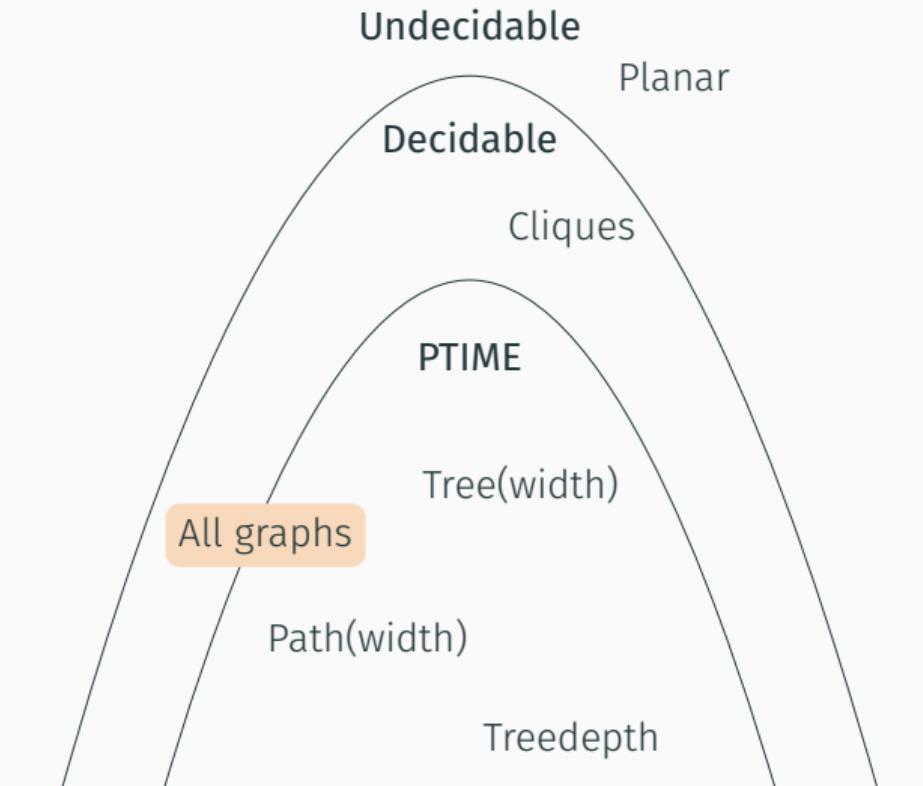
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An Algorithmic Meta Theorem

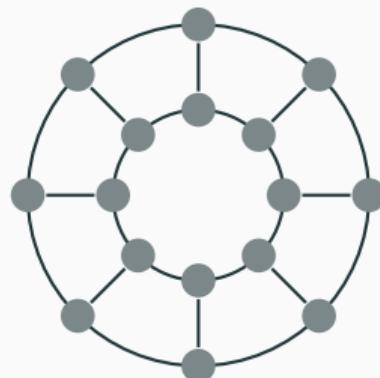
Theorem (S. (2024))

For every *recognisable graph class \mathcal{F} of bounded treewidth*, $\text{HOMIND}(\mathcal{F})$ is in **coRP**.

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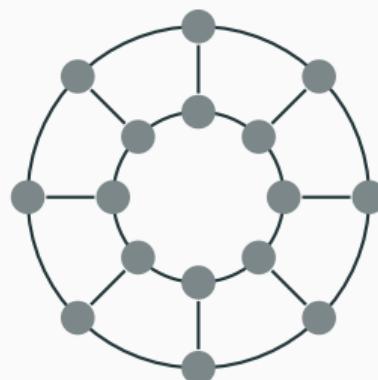


Minor-closed and **bounded**
treewidth

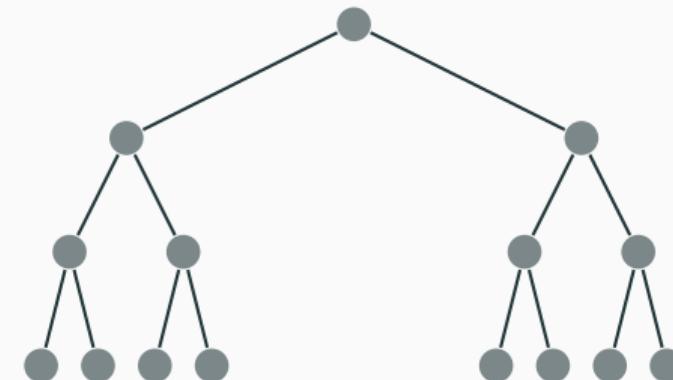
An Algorithmic Meta Theorem

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Minor-closed and bounded
treewidth



CMSO₂-definable and bounded treewidth

Courcelle (1990); Bojańczyk and Pilipczuk (2016)

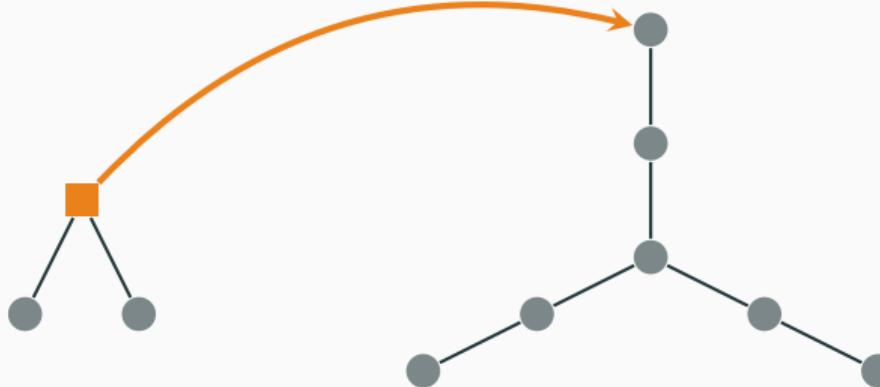
Labelled Graphs and Homomorphism Vectors



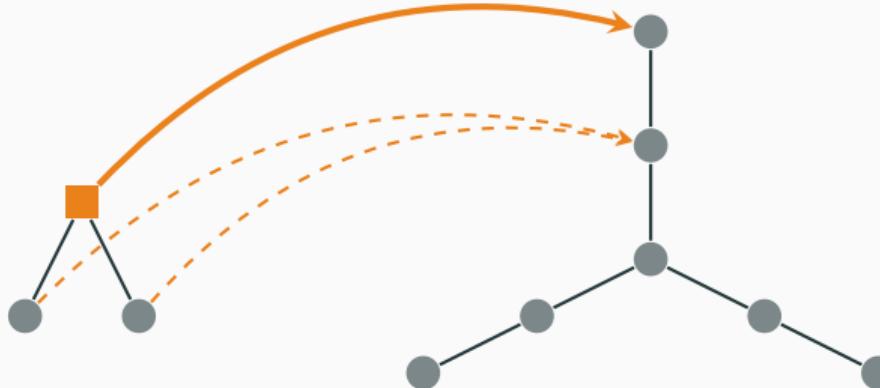
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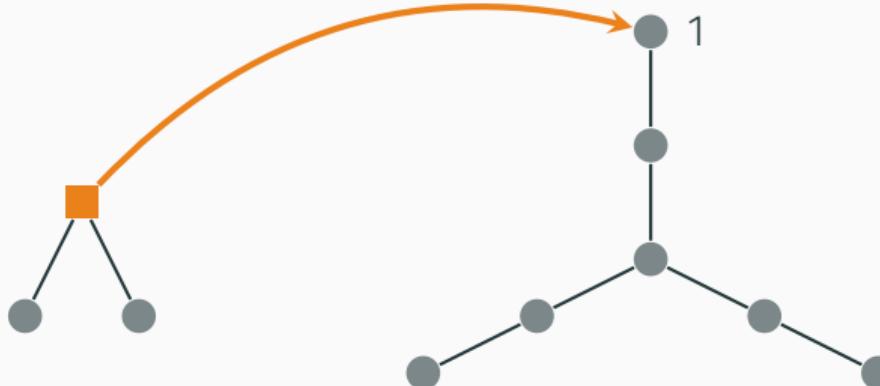
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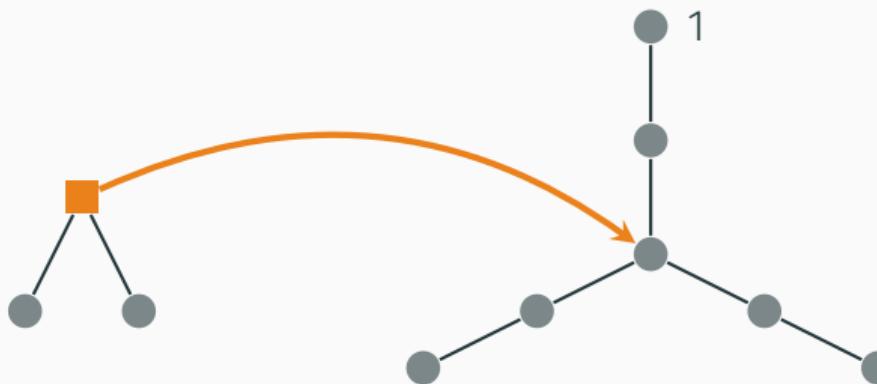
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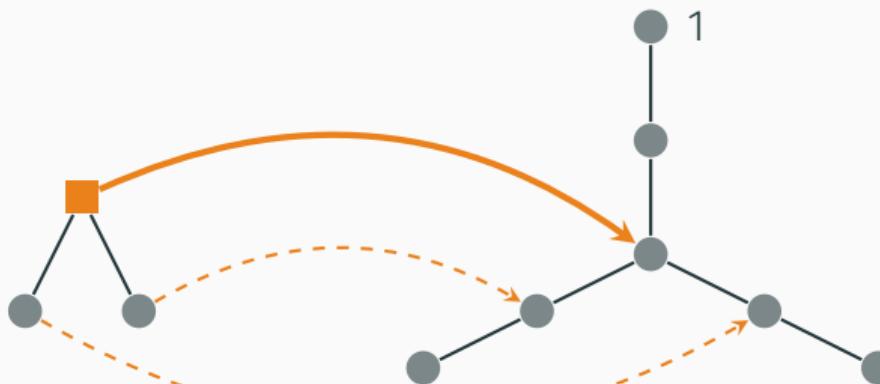
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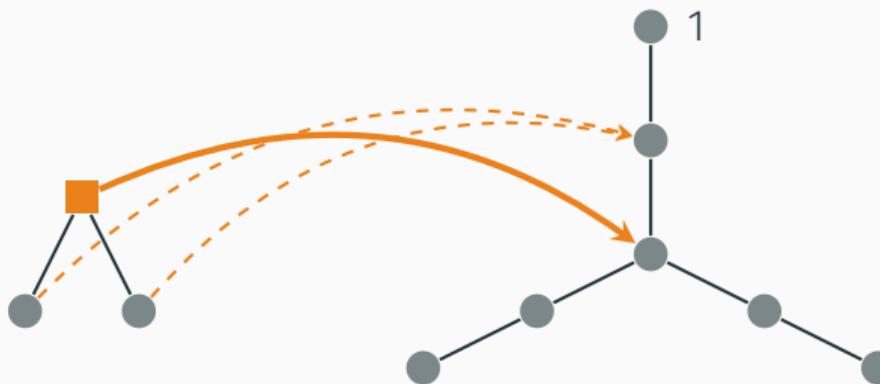
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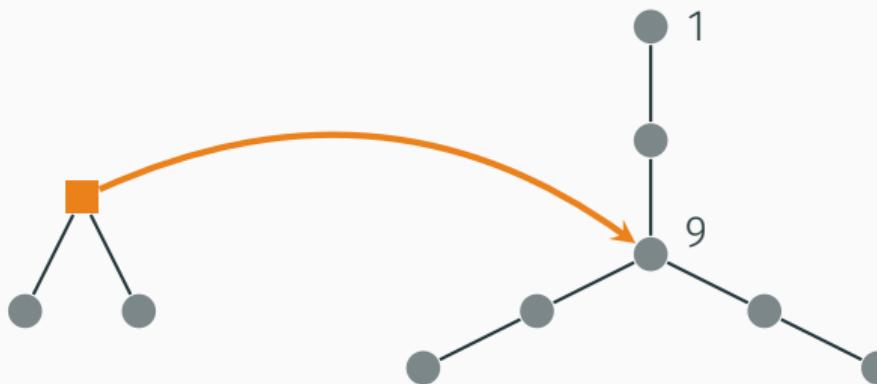
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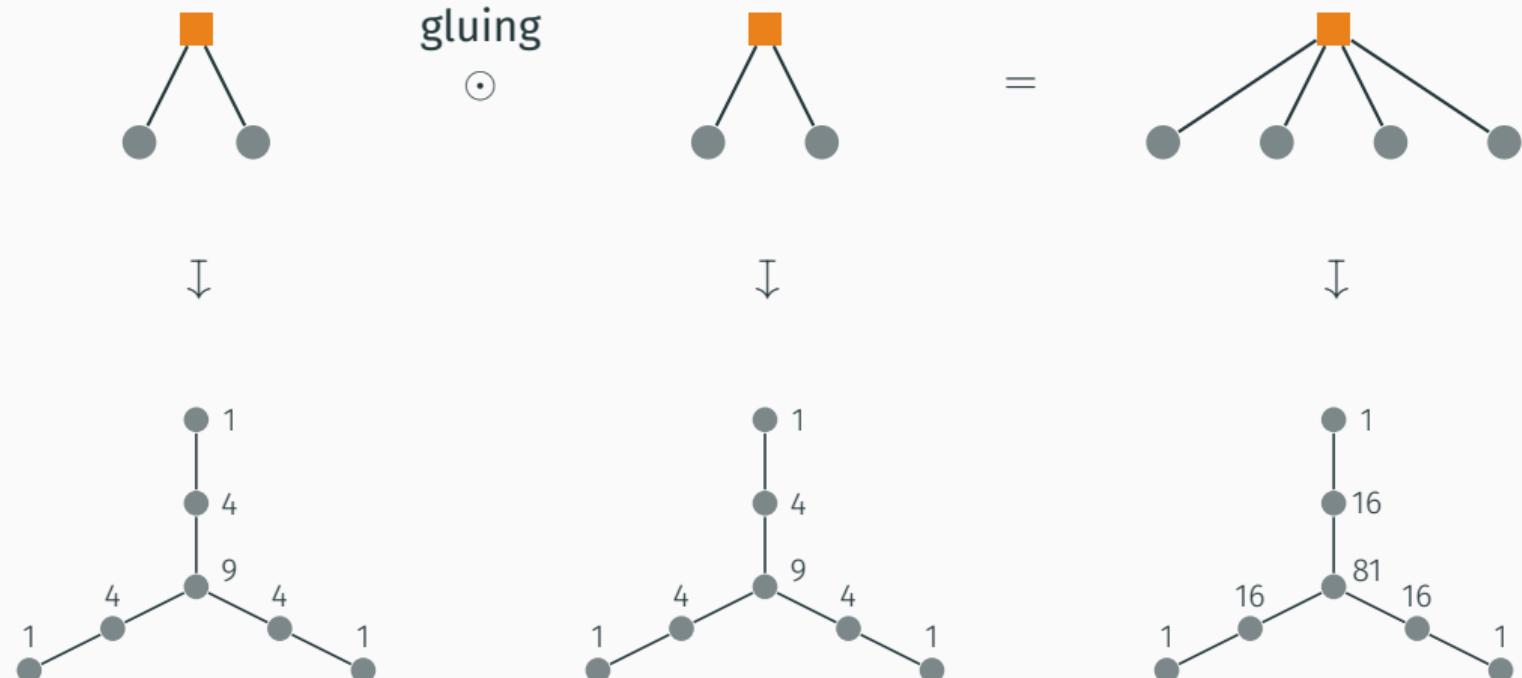
$$\mathcal{F} \longrightarrow \mathbb{R}^{V(G)}$$



Combinatorial and Algebraic Operations: Gluing and Schur Product



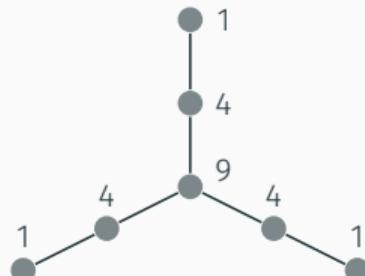
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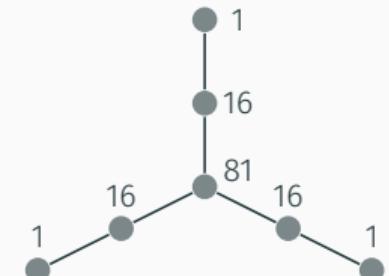
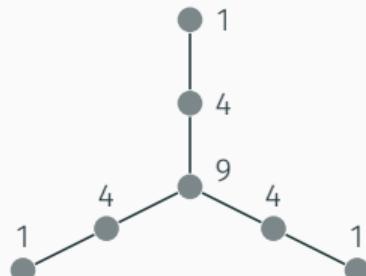
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gluing



Schur
product



Homomorphism Indistinguishability over Trees

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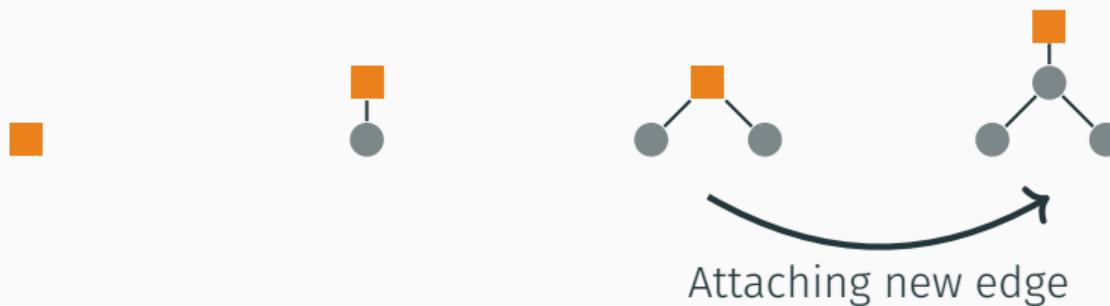
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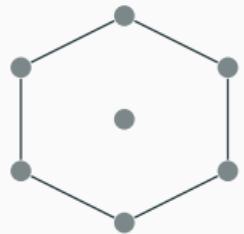


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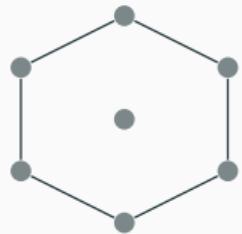


Homomorphism Indistinguishability over Trees





Space of homomorphism vectors of labelled trees



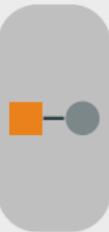
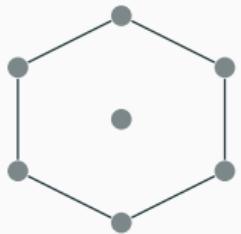
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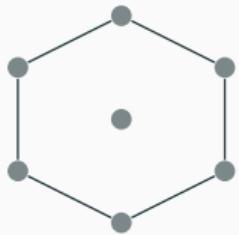


Attaching new edge



Space of homomorphism vectors of labelled trees





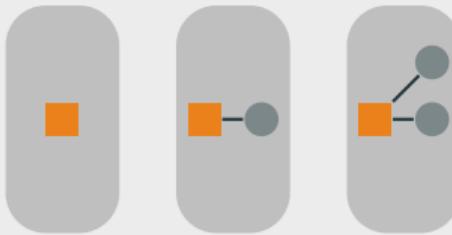
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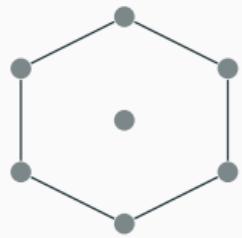


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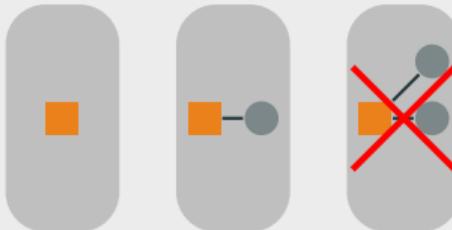
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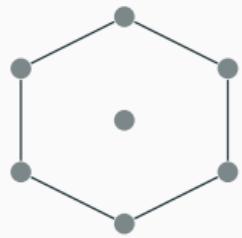


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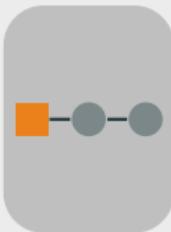
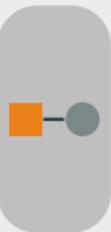
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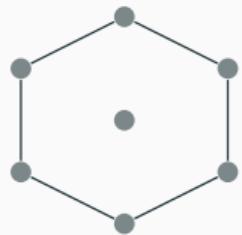


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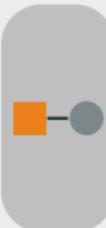
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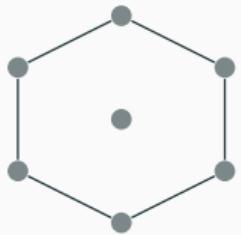


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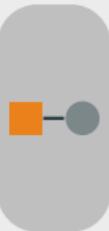
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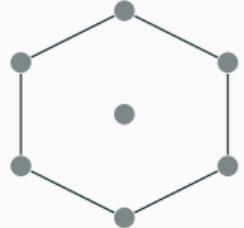


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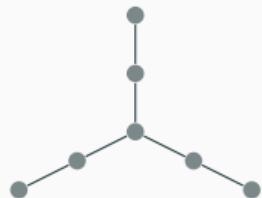
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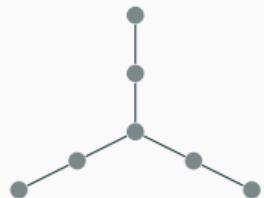
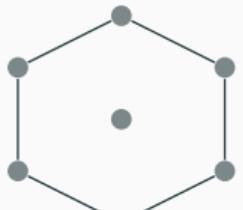
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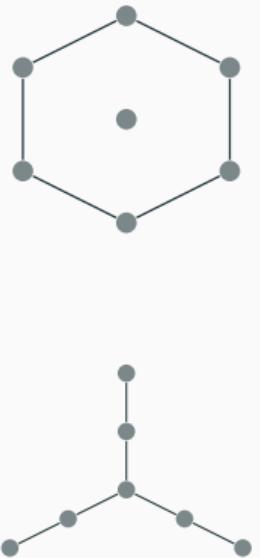


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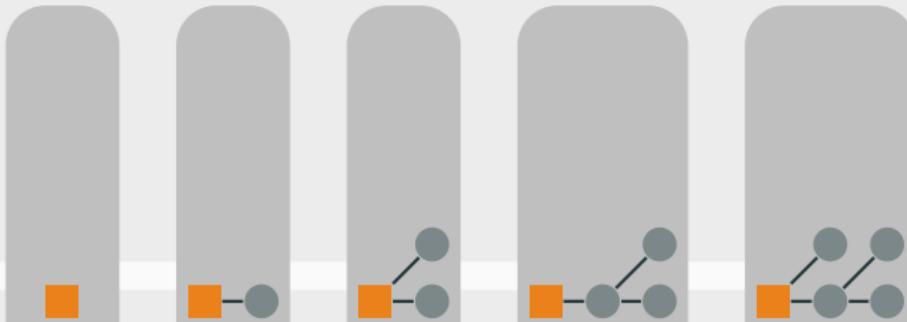


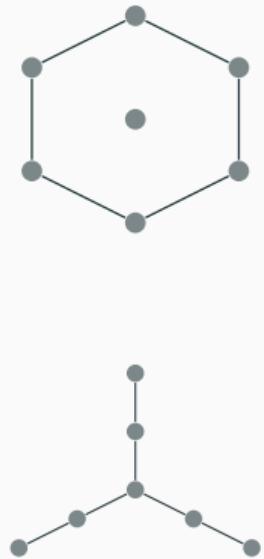
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Space of homomorphism vectors of labelled trees





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Space of homomorphism vectors of labelled trees

7



12



24



48



192



54



240

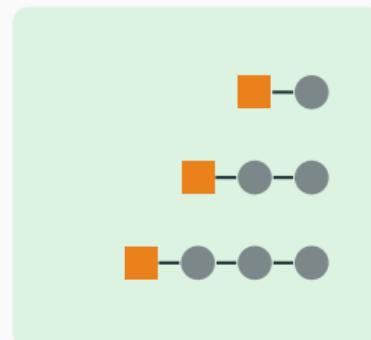


Let \mathcal{F} be a graph class. Let $F_1 \sim_{\mathcal{F}} F_2$ iff for all labelled graphs K

$$\text{unlabel}(K \odot F_1) \in \mathcal{F} \iff \text{unlabel}(K \odot F_2) \in \mathcal{F}.$$

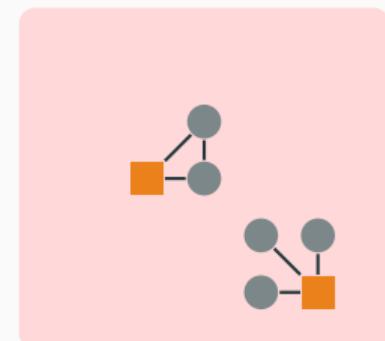
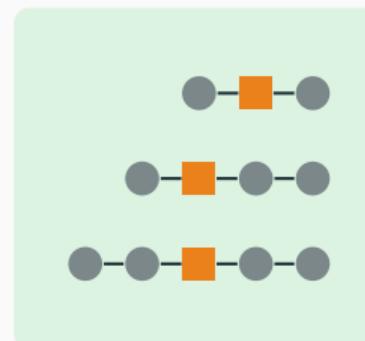
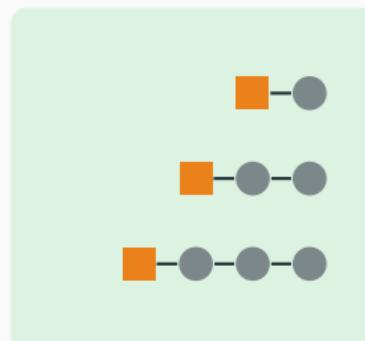
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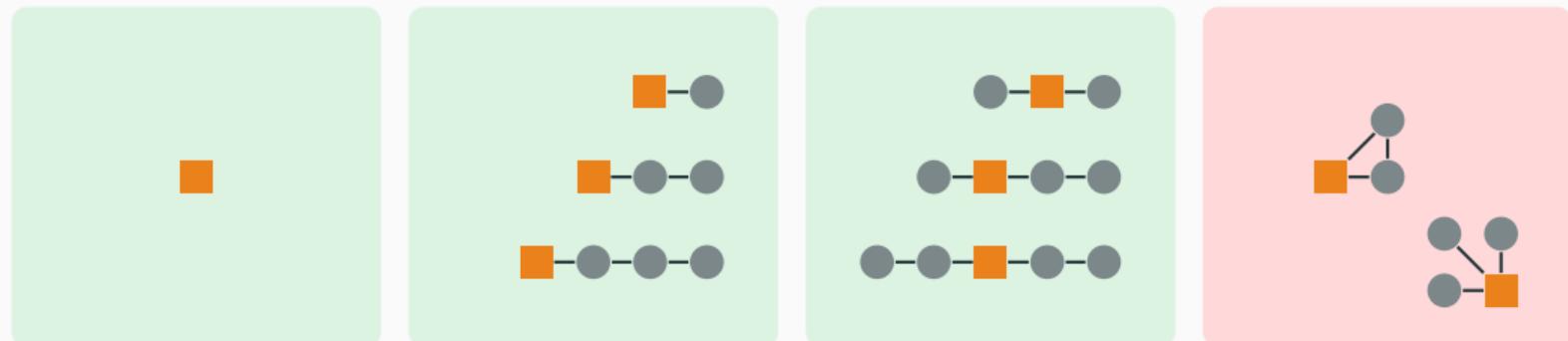
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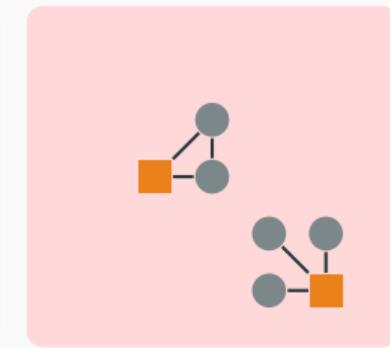
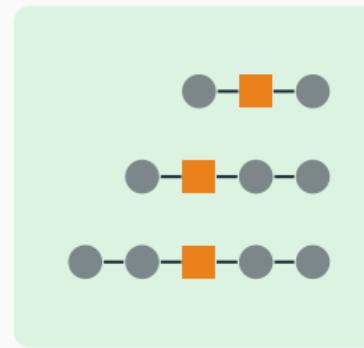
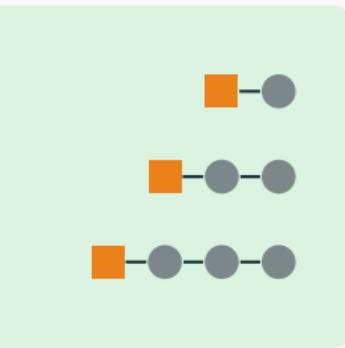
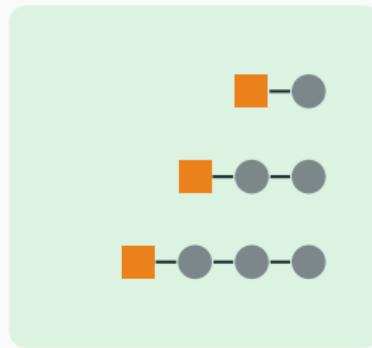
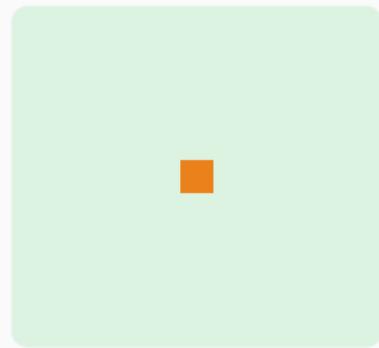


Definition

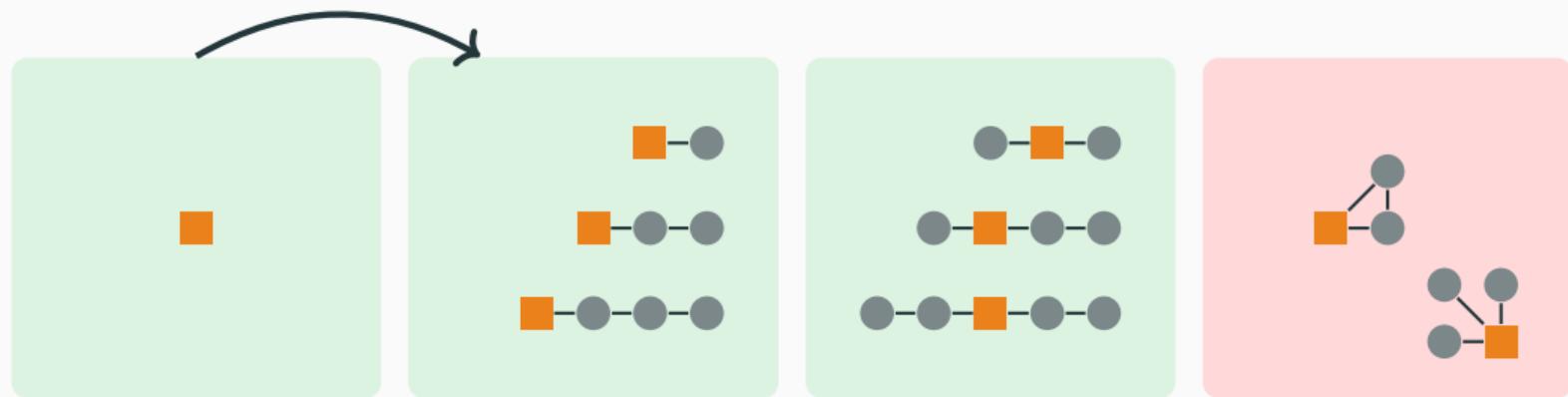
\mathcal{F} is *recognisable* if $\sim_{\mathcal{F}}$ has finitely many equivalence classes.

Attaching a new edge

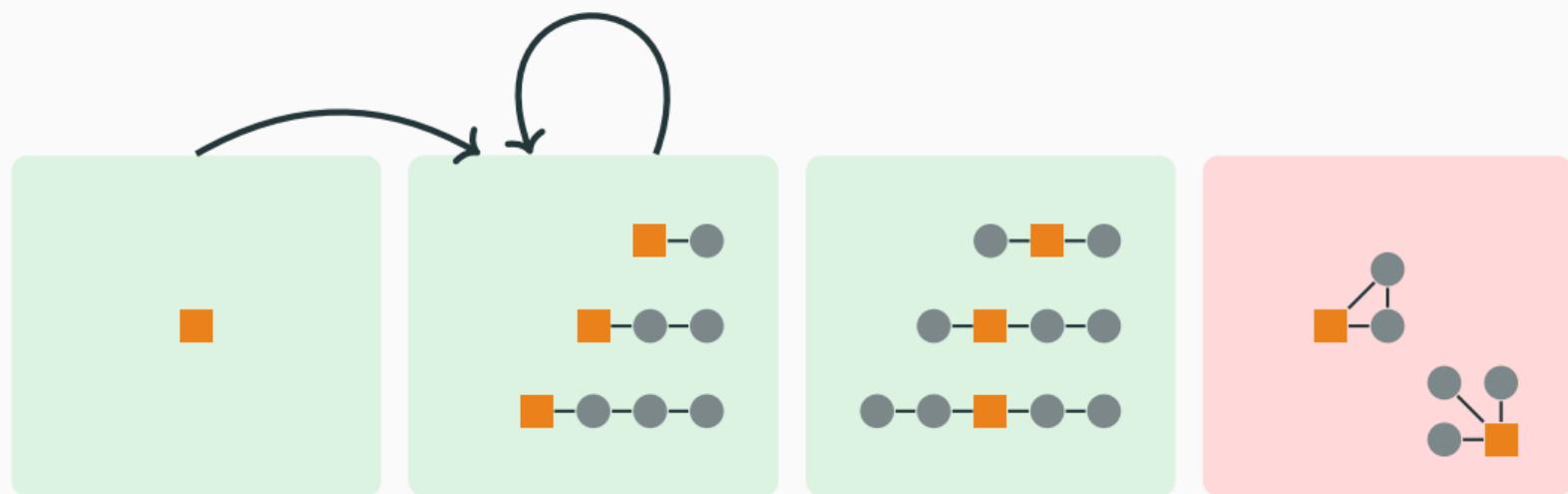
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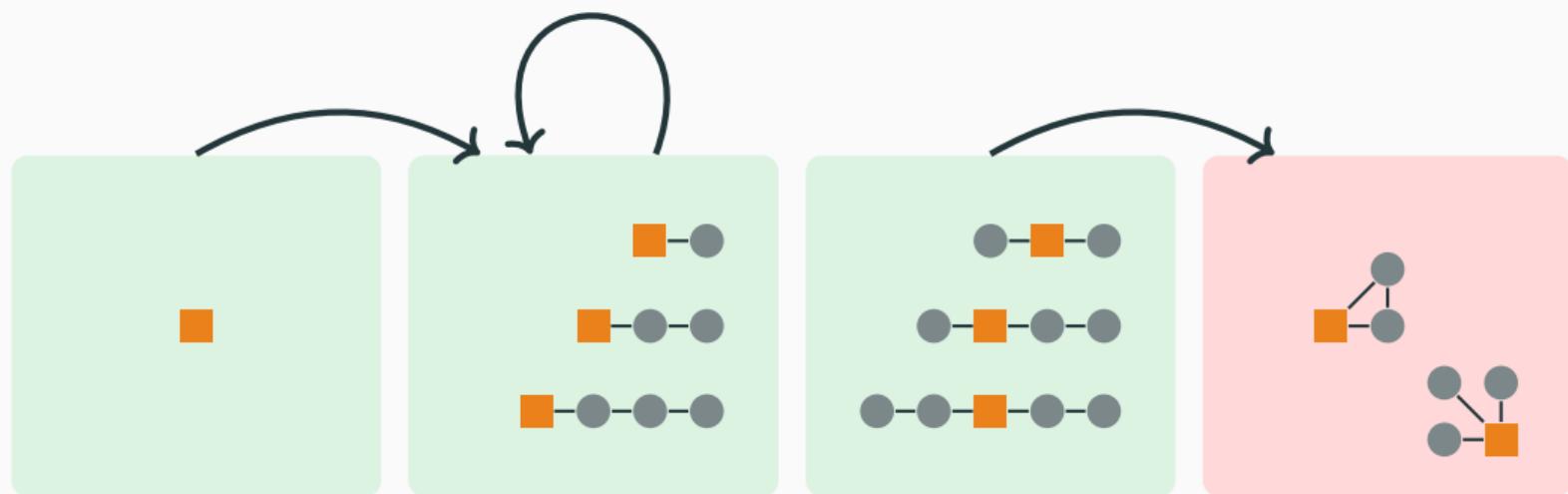
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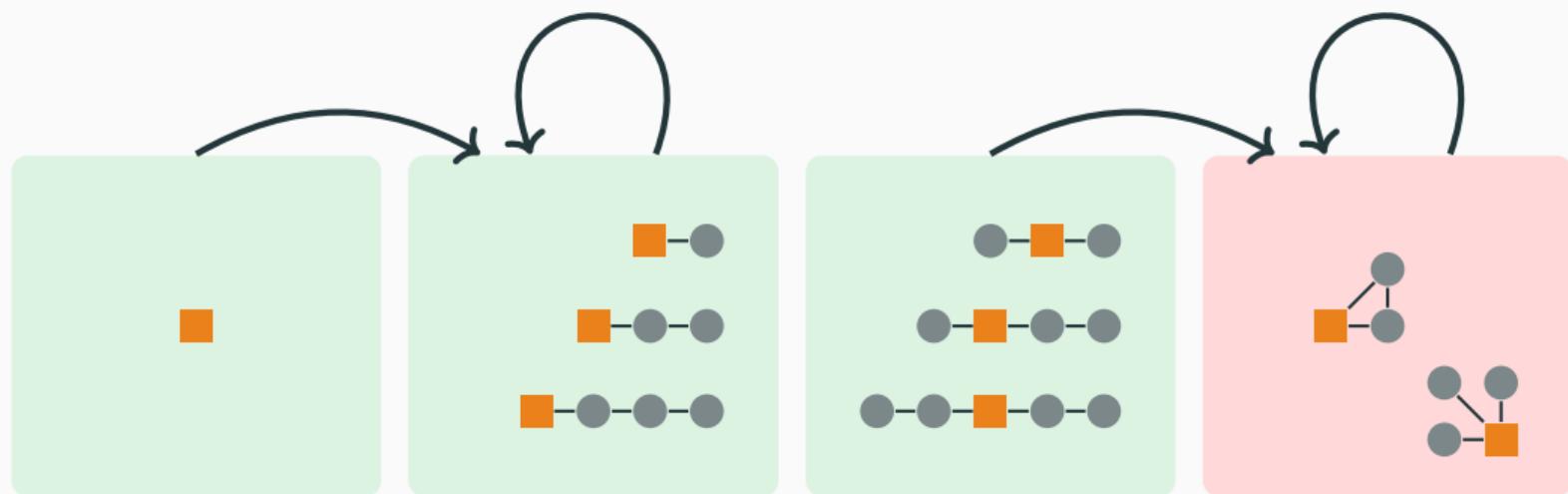
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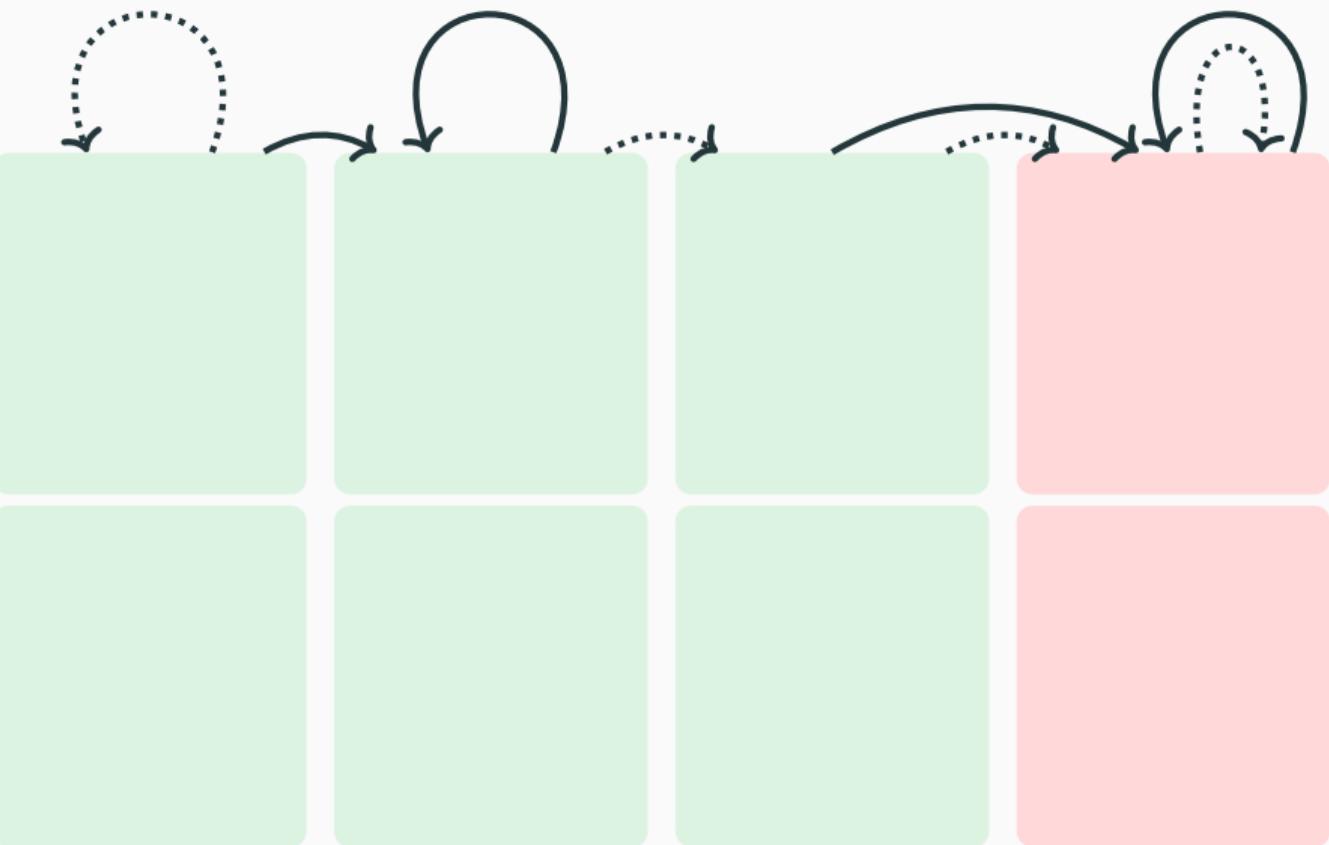
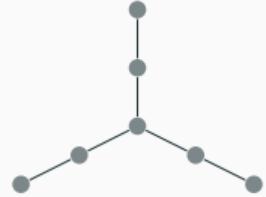
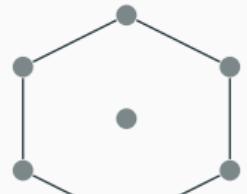


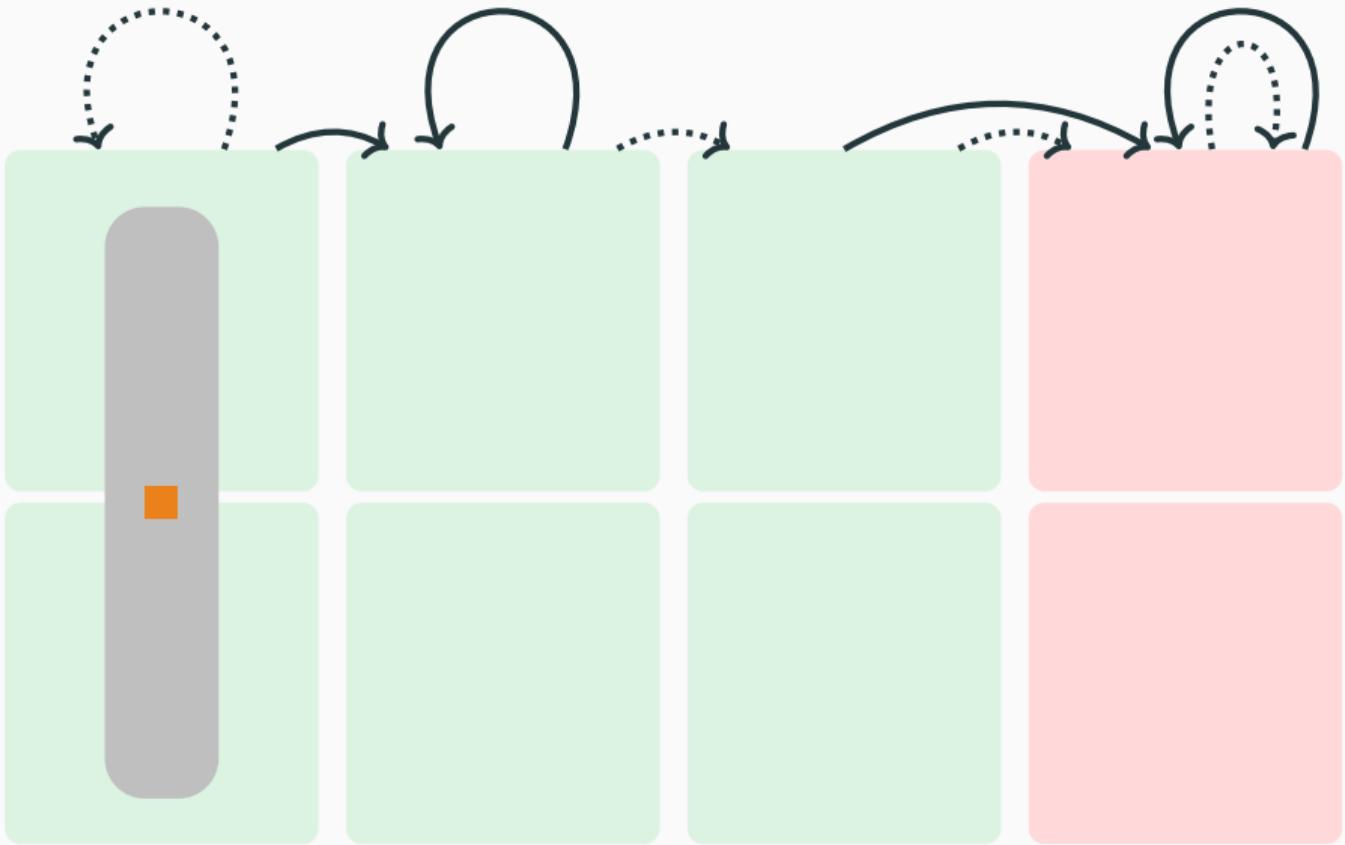
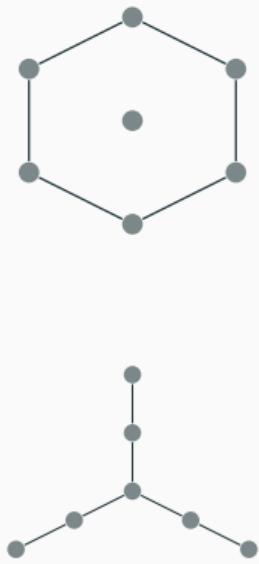
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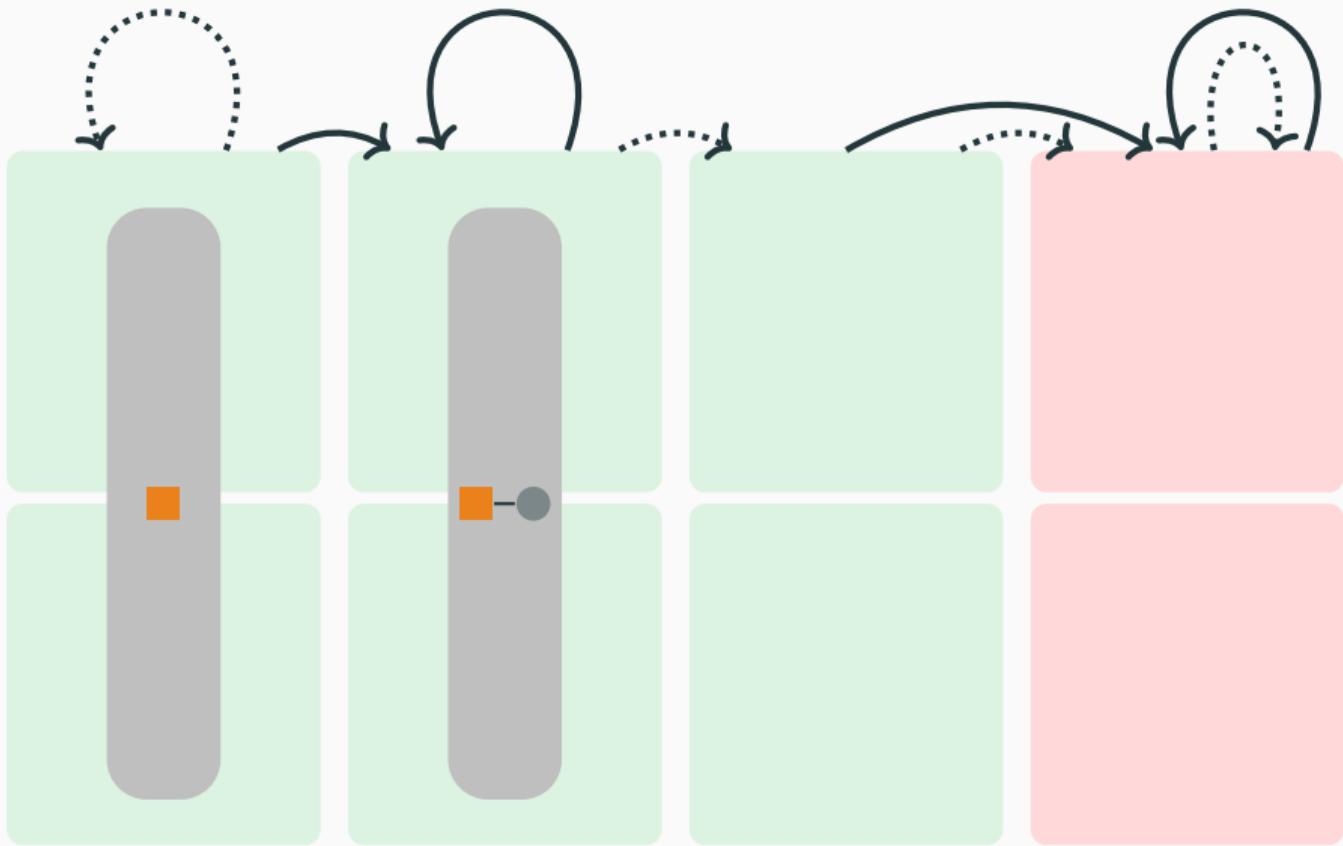
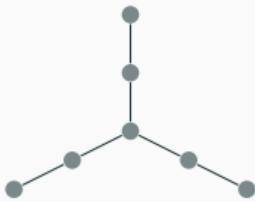
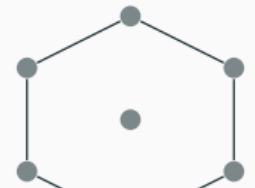


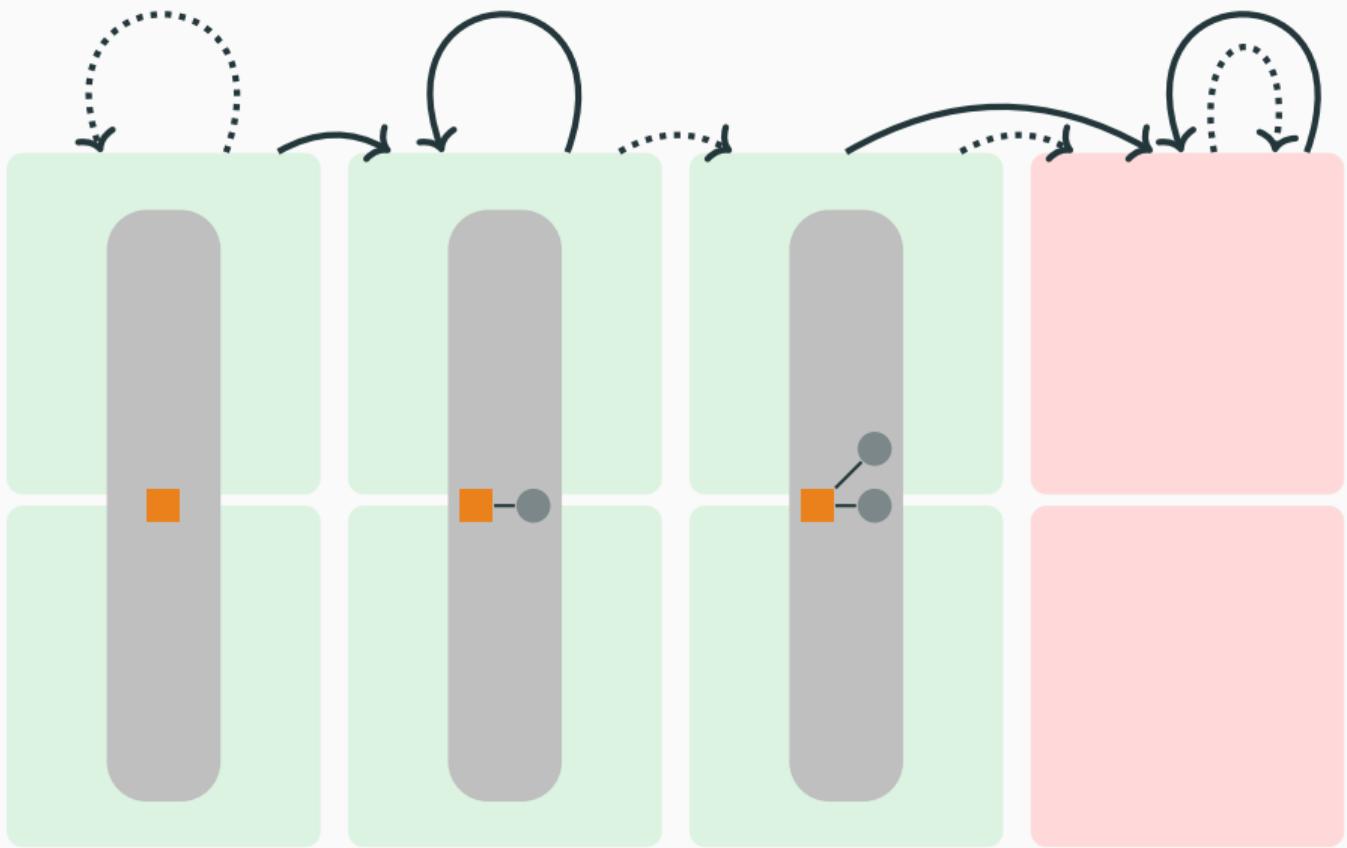
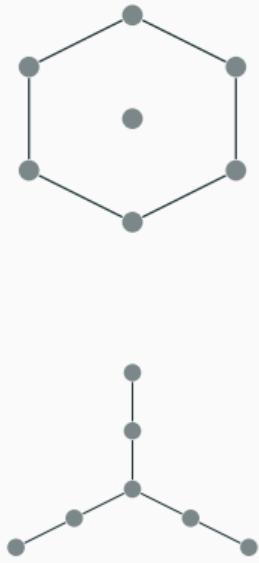
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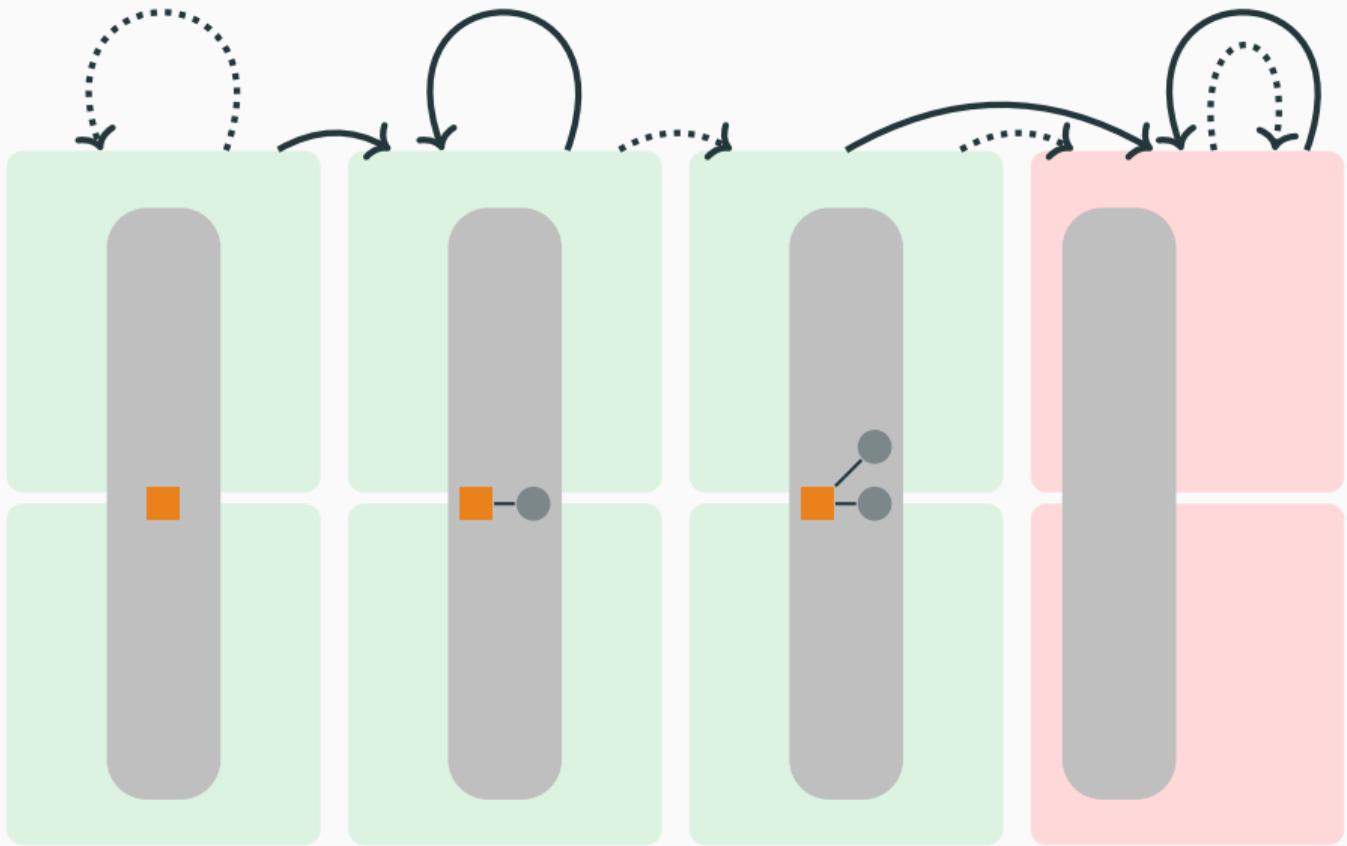
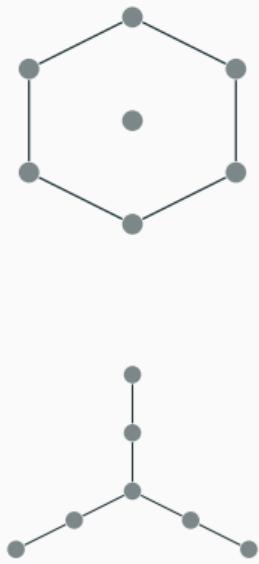


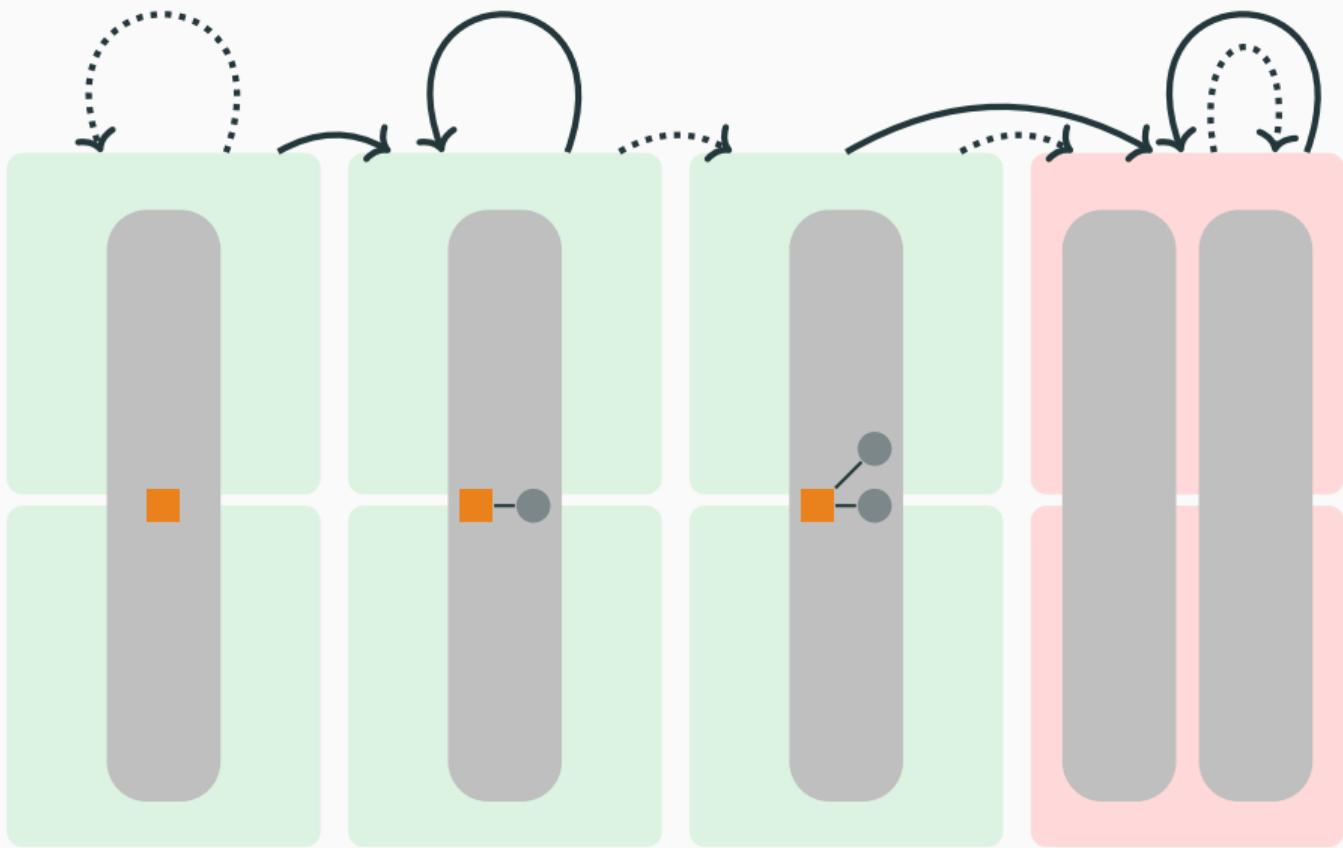
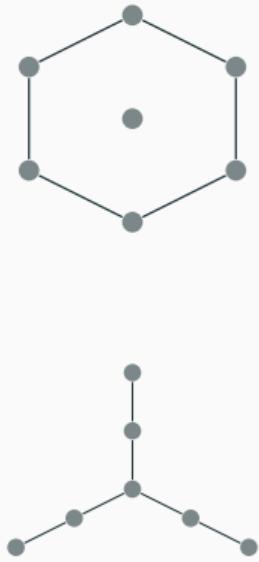


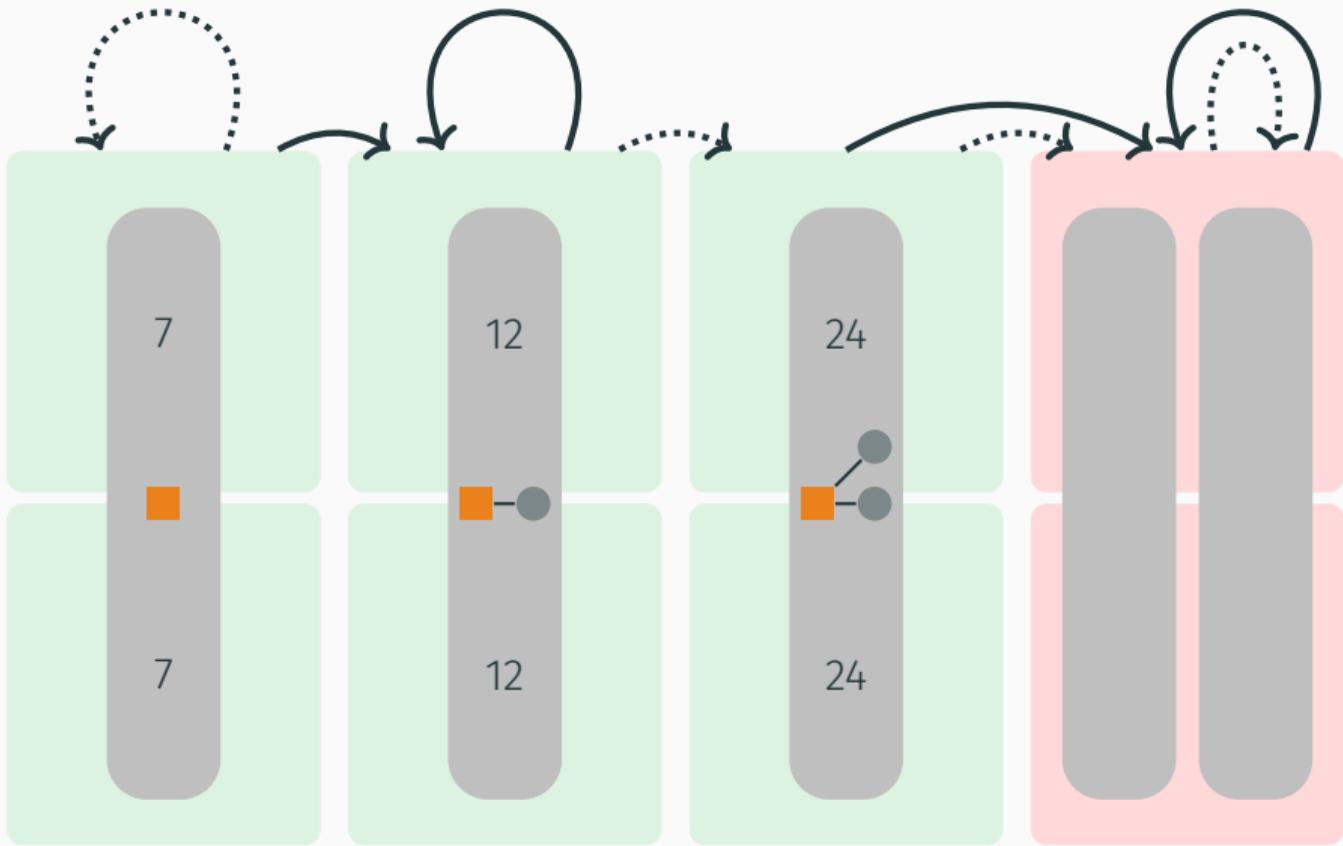
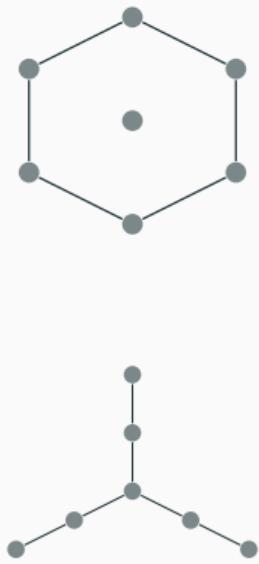










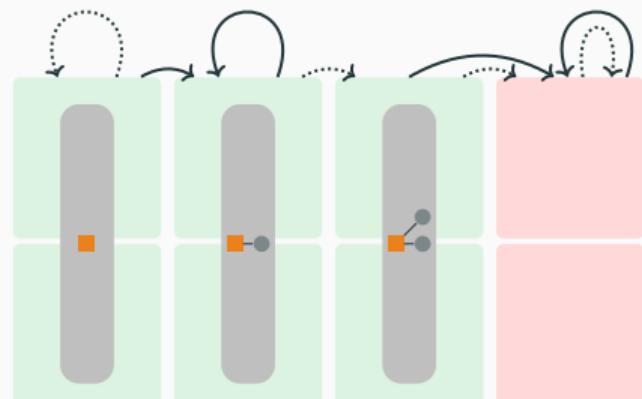


Conclusion

Theorem (S. (2024))

For every *recognisable graph class \mathcal{F} of bounded treewidth*, $\text{HOMIND}(\mathcal{F})$ is in **coRP**.

- Randomised fixed-parameter algorithm for CMSO_2 -sentence specifying \mathcal{F} .



A conjecture

Conjecture

For a minor-closed \mathcal{F} , either

\mathcal{F} contains *all graphs* and $\text{HOMIND}(\mathcal{F})$ is *Graph Isomorphism*,

A conjecture

Conjecture

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\mathcal{F} contains *all graphs* and $\text{HOMIND}(\mathcal{F})$ is *Graph Isomorphism*,

\mathcal{F} has *bounded treewidth* and $\text{HOMIND}(\mathcal{F})$ is in *coRP*, or

A conjecture

Conjecture

For a minor-closed \mathcal{F} , either

\mathcal{F} contains *all graphs* and $\text{HOMIND}(\mathcal{F})$ is *Graph Isomorphism*,

\mathcal{F} has *bounded treewidth* and $\text{HOMIND}(\mathcal{F})$ is in *coRP*, or

\mathcal{F} has *unbounded treewidth* and $\text{HOMIND}(\mathcal{F})$ is *undecidable*.

Parametrised HomIND

HomIND

Input Graphs G and H , a CSMO₂-sentence φ , an integer $k \in \mathbb{N}$

Parameter $|\varphi| + k$

Decide Are G and H homomorphism indistinguishable over

$$\mathcal{F}_{\varphi,k} := \{F \mid \text{tw } F \leq k, F \models \varphi\}?$$

Parametrised HOMIND

HOMIND

Input Graphs G and H , a CSMO₂-sentence φ , an integer $k \in \mathbb{N}$

Parameter $|\varphi| + k$

Decide Are G and H homomorphism indistinguishable over

$$\mathcal{F}_{\varphi,k} := \{F \mid \text{tw } F \leq k, F \models \varphi\}?$$

Theorem (S. (2024))

There is a randomised algorithm for HOMIND running in time $f(k, |\varphi|)n^{O(k)}$.

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Bibliography iii

Title Picture: "Bicycle race scene. A peloton of six cyclists crosses the finish line in front of a crowded grandstand, observed by a referee." (1895) by Calvert Lithographic Co., Detroit, Michigan, Public Domain, via Wikimedia Commons. https://commons.wikimedia.org/wiki/File:Bicycle_race_scene,_1895.jpg