

A vintage painting depicting a bicycle race. In the foreground, several cyclists are shown in motion, wearing colorful jerseys and shorts. A jockey in a brown suit and hat is seen from the back, looking at a clipboard. In the background, a large crowd of spectators in early 20th-century attire is gathered behind a white picket fence, watching the race. The scene is set outdoors with a clear sky and some trees in the distance.

An Algorithmic Meta Theorem for Homomorphism Indistinguishability

MFCS 2024, 29 August 2024

Tim Seppelt



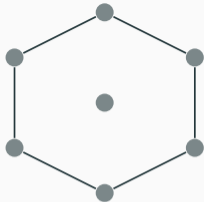
Research Training Group –
Uncertainty and Randomness
in Algorithms, Verification,
and Logic

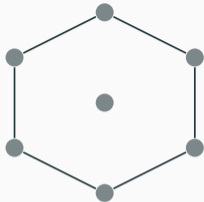
RWTHAACHEN
UNIVERSITY

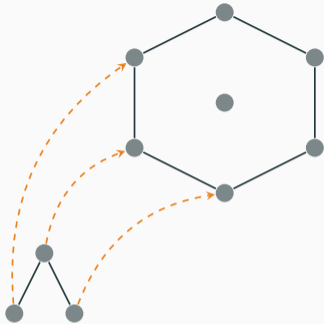


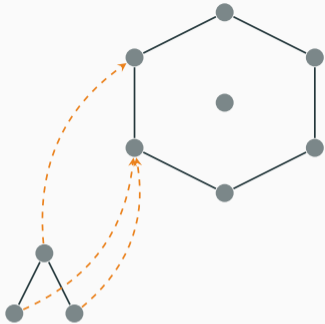
Chair for Logic
and Theory of
Discrete Systems

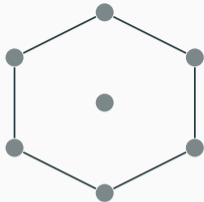
RWTHAACHEN
UNIVERSITY



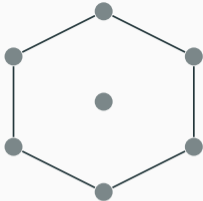








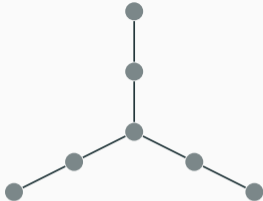
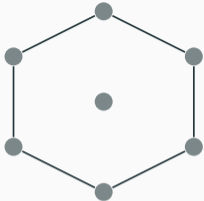
24



24



36

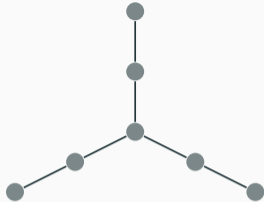
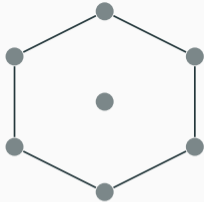


24

24

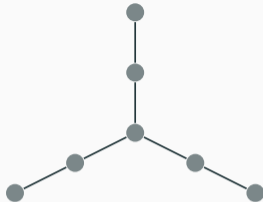
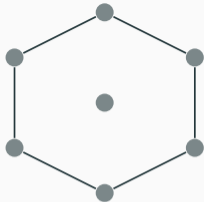
36

36



24
36

24
36





24

36



24

36

The graphs  and  are homomorphism indistinguishable over $\left\{ \begin{array}{c} \text{triangle with one vertex removed} \\ \text{square cycle} \end{array} \right\}$.

First Motivation: Characterisations and Connections

graph class \mathcal{F} relation $\equiv_{\mathcal{F}}$
 \mathcal{G} isomorphism

Lovász (1967)

First Motivation: Characterisations and Connections

graph class \mathcal{F}	relation $\equiv_{\mathcal{F}}$	
\mathcal{G}	isomorphism	Lovász (1967)
\mathcal{C}	cospectral adjacency matrices	Folklore

First Motivation: Characterisations and Connections

graph class \mathcal{F}	relation $\equiv_{\mathcal{F}}$	
\mathcal{G}	isomorphism	Lovász (1967)
\mathcal{C}	cospectral adjacency matrices	Folklore
\mathcal{P}	quantum isomorphism	Mančinska and Roberson (2020)

First Motivation: Characterisations and Connections

graph class \mathcal{F}	relation $\equiv_{\mathcal{F}}$	
\mathcal{G}	isomorphism	Lovász (1967)
\mathcal{C}	cospectral adjacency matrices	Folklore
\mathcal{P}	quantum isomorphism	Mančinska and Roberson (2020)
\mathcal{TW}_k	C^{k+1} -equivalence	Dvořák (2010); Dell, Grohe, and Rattan (2018)

First Motivation: Characterisations and Connections

graph class \mathcal{F}	relation $\equiv_{\mathcal{F}}$	
\mathcal{G}	isomorphism	Lovász (1967)
\mathcal{C}	cospectral adjacency matrices	Folklore
\mathcal{P}	quantum isomorphism	Mančinska and Roberson (2020)
\mathcal{TW}_k	C^{k+1} -equivalence	Dvořák (2010); Dell, Grohe, and Rattan (2018)
	Weisfeiler–Leman algorithm	Cai, Fürer, and Immerman (1992)

First Motivation: Characterisations and Connections

graph class \mathcal{F}	relation $\equiv_{\mathcal{F}}$	
\mathcal{G}	isomorphism	Lovász (1967)
\mathcal{C}	cospectral adjacency matrices	Folklore
\mathcal{P}	quantum isomorphism	Mančinska and Roberson (2020)
\mathcal{TW}_k	C^{k+1} -equivalence	Dvořák (2010); Dell, Grohe, and Rattan (2018)
	Weisfeiler–Leman algorithm	Cai, Fürer, and Immerman (1992)
	Sherali–Adams hierarchy	Atserias and Maneva (2012); Grohe and Otto (2015)

First Motivation: Characterisations and Connections

graph class \mathcal{F}	relation $\equiv_{\mathcal{F}}$	
\mathcal{G}	isomorphism	Lovász (1967)
\mathcal{C}	cospectral adjacency matrices	Folklore
\mathcal{P}	quantum isomorphism	Mančinska and Roberson (2020)
\mathcal{TW}_k	C^{k+1} -equivalence	Dvořák (2010); Dell, Grohe, and Rattan (2018)
	Weisfeiler–Leman algorithm	Cai, Fürer, and Immerman (1992)
	Sherali–Adams hierarchy	Atserias and Maneva (2012); Grohe and Otto (2015)
	\mathbb{P}_k -coKleisli-isomorphism	Dawar, Jakl, and Reggio (2021)

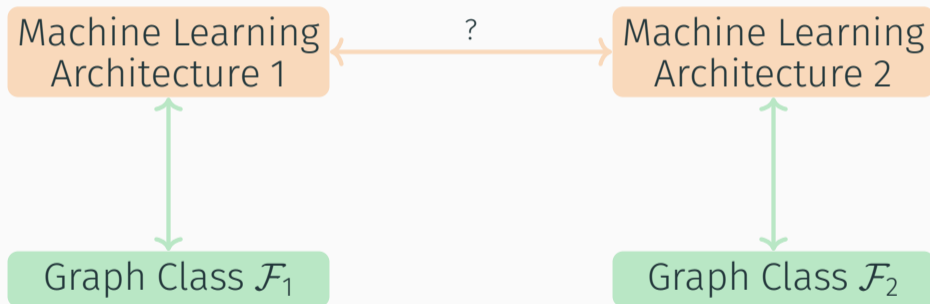
First Motivation: Characterisations and Connections

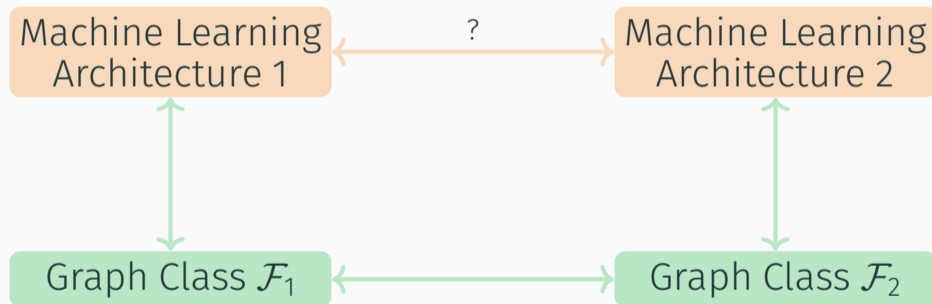
graph class \mathcal{F}	relation $\equiv_{\mathcal{F}}$	
\mathcal{G}	isomorphism	Lovász (1967)
\mathcal{C}	cospectral adjacency matrices	Folklore
\mathcal{P}	quantum isomorphism	Mančinska and Roberson (2020)
\mathcal{TW}_k	C^{k+1} -equivalence	Dvořák (2010); Dell, Grohe, and Rattan (2018)
	Weisfeiler–Leman algorithm	Cai, Fürer, and Immerman (1992)
	Sherali–Adams hierarchy	Atserias and Maneva (2012); Grohe and Otto (2015)
	\mathbb{P}_k -coKleisli-isomorphism	Dawar, Jakl, and Reggio (2021)
	Graph Neural Networks	Xu et al. (2018); Morris et al. (2019)
...	...	

Machine Learning
Architecture 1

?

Machine Learning
Architecture 2





Complexity of Homomorphism Indistinguishability

HOMIND(\mathcal{F})

Input Graphs G and H .

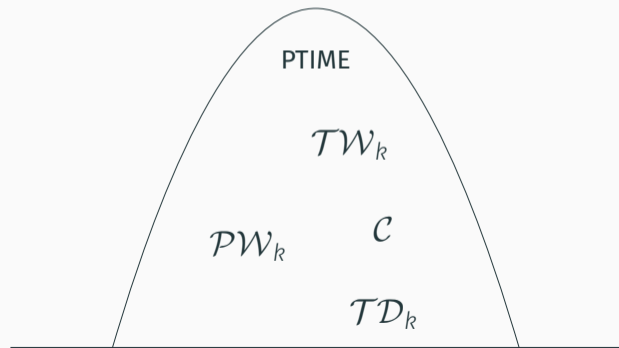
Decide $G \equiv_{\mathcal{F}} H$.

Complexity of Homomorphism Indistinguishability

$\text{HOMIND}(\mathcal{F})$

Input Graphs G and H .

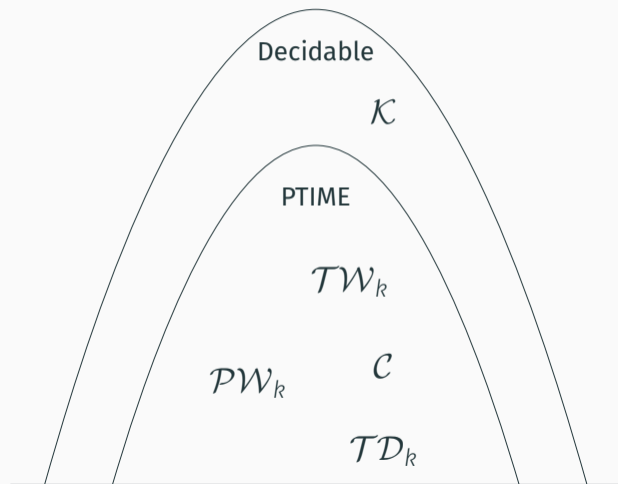
Decide $G \equiv_{\mathcal{F}} H$.



$\text{HOMIND}(\mathcal{F})$

Input Graphs G and H .

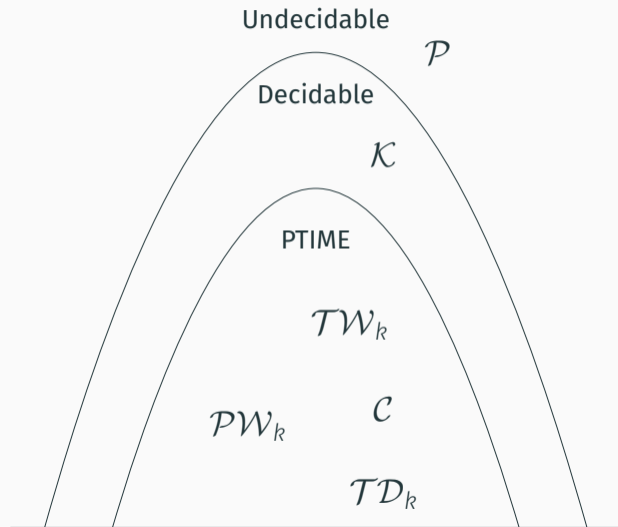
Decide $G \equiv_{\mathcal{F}} H$.



$\text{HOMIND}(\mathcal{F})$

Input Graphs G and H .

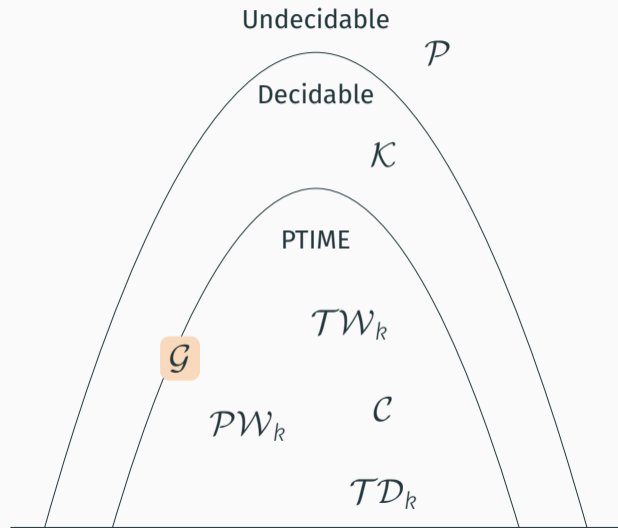
Decide $G \equiv_{\mathcal{F}} H$.



$\text{HOMIND}(\mathcal{F})$

Input Graphs G and H .

Decide $G \equiv_{\mathcal{F}} H$.



An Algorithmic Meta Theorem

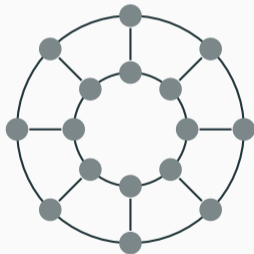
Theorem

For every *recognisable* graph class \mathcal{F} of *bounded treewidth*, $\text{HOMIND}(\mathcal{F})$ is in **coRP**.

An Algorithmic Meta Theorem

Theorem

For every *recognisable* graph class \mathcal{F} of *bounded treewidth*, $\text{HOMIND}(\mathcal{F})$ is in *coRP*.

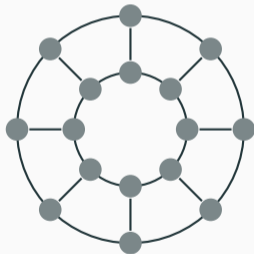


Minor-closed and bounded
treewidth

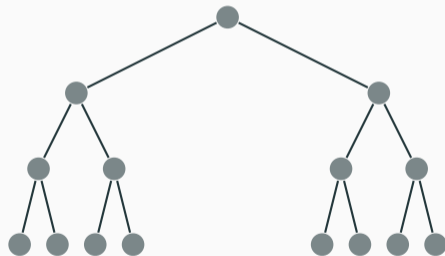
An Algorithmic Meta Theorem

Theorem

For every *recognisable* graph class \mathcal{F} of *bounded treewidth*, $\text{HOMIND}(\mathcal{F})$ is in **coRP**.



Minor-closed and bounded
treewidth



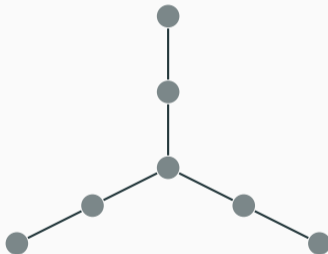
CMSO₂-definable and bounded treewidth

Courcelle (1990); Bojańczyk and Pilipczuk (2016)

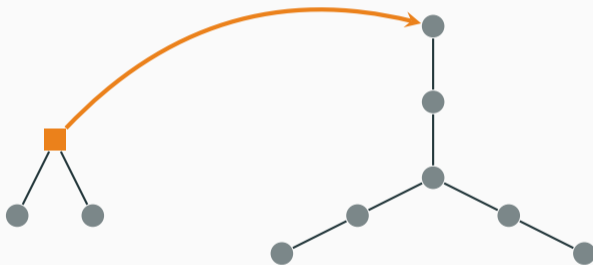
Labelled Graphs and Homomorphism Vectors



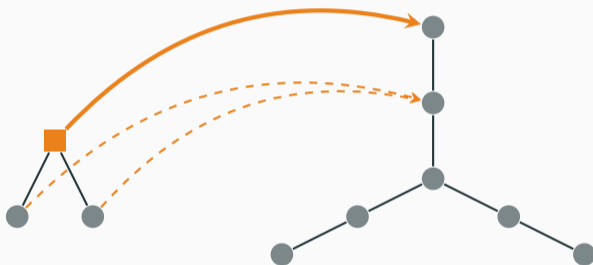
Labelled Graphs and Homomorphism Vectors



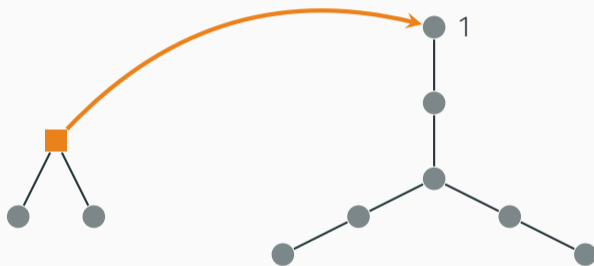
Labelled Graphs and Homomorphism Vectors



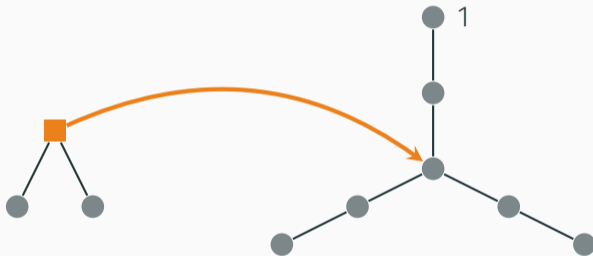
Labelled Graphs and Homomorphism Vectors



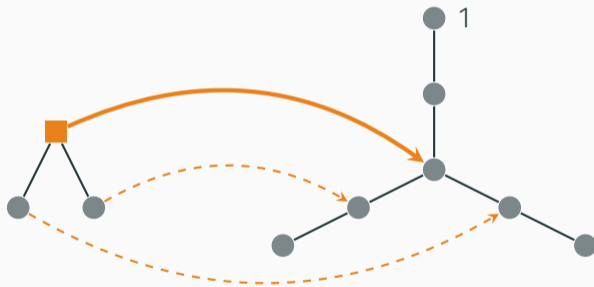
Labelled Graphs and Homomorphism Vectors



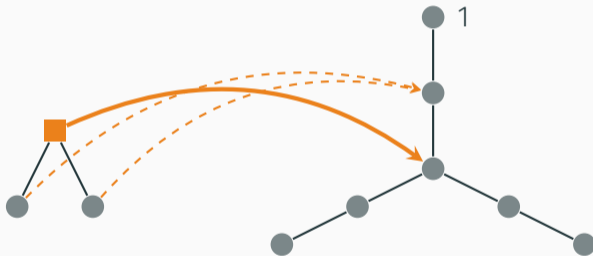
Labelled Graphs and Homomorphism Vectors



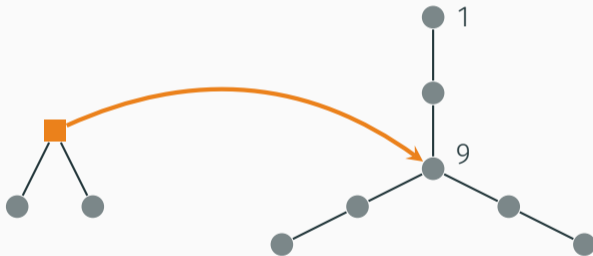
Labelled Graphs and Homomorphism Vectors



Labelled Graphs and Homomorphism Vectors

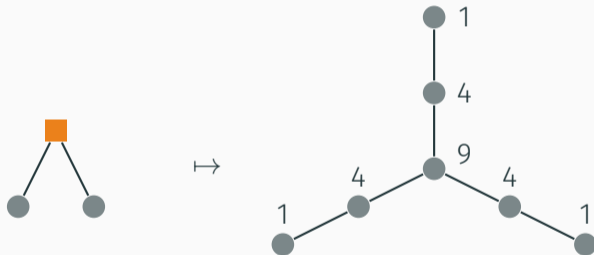


Labelled Graphs and Homomorphism Vectors



Labelled Graphs and Homomorphism Vectors

$$\mathcal{F} \longrightarrow \mathbb{R}^{V(G)}$$



Combinatorial and Algebraic Operations: Gluing and Schur Product



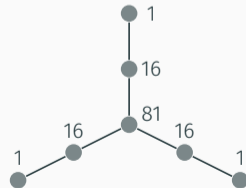
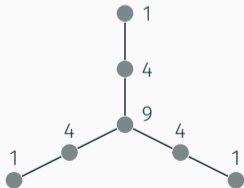
Combinatorial and Algebraic Operations: Gluing and Schur Product



gluing



=



Combinatorial and Algebraic Operations: Gluing and Schur Product



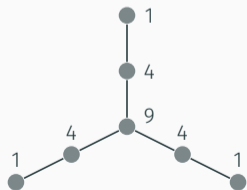
gluing



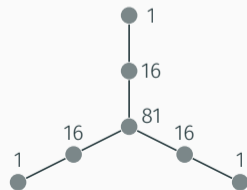
=



Schur
product



=



Homomorphism Indistinguishability over Trees

Homomorphism Indistinguishability over Trees



Homomorphism Indistinguishability over Trees

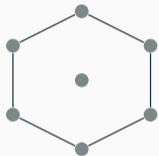


Homomorphism Indistinguishability over Trees

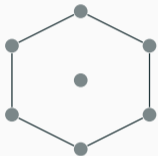


Homomorphism Indistinguishability over Trees





Space of homomorphism vectors of labelled trees



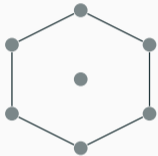
Space of homomorphism vectors of labelled trees

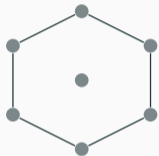


Attaching new edge



Space of homomorphism vectors of labelled trees





Attaching new edge

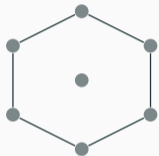


Gluing



Space of homomorphism vectors of labelled trees





Attaching new edge

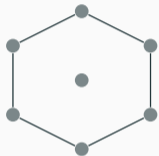


Gluing



Space of homomorphism vectors of labelled trees





Attaching new edge



Gluing



Space of homomorphism vectors of labelled trees

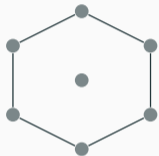


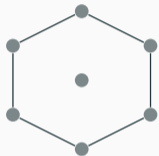
Attaching new edge

Gluing



Space of homomorphism vectors of labelled trees





Attaching new edge



Gluing



Space of homomorphism vectors of labelled trees



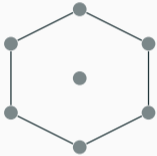
Attaching new edge



Gluing



Space of homomorphism vectors of labelled trees

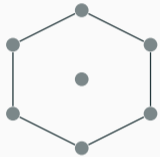


Attaching new edge

Gluing



Space of homomorphism vectors of labelled trees

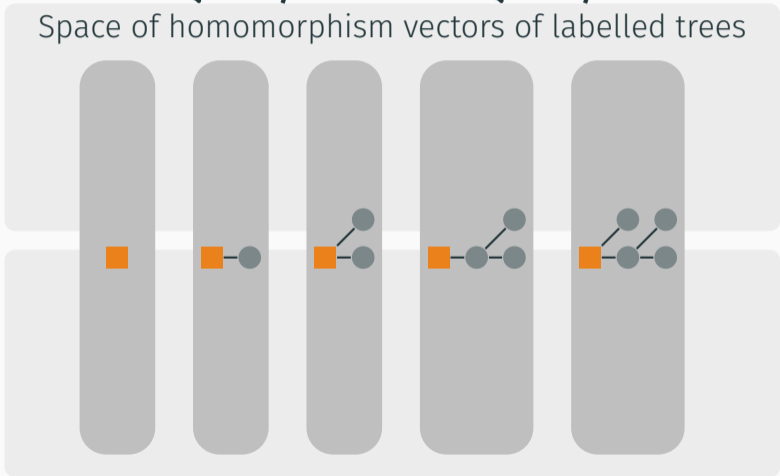


Attaching new edge

Gluing



Space of homomorphism vectors of labelled trees

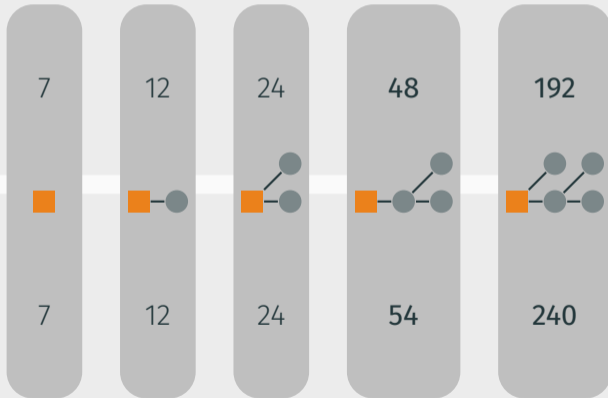
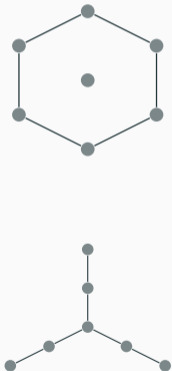


Attaching new edge

Gluing



Space of homomorphism vectors of labelled trees

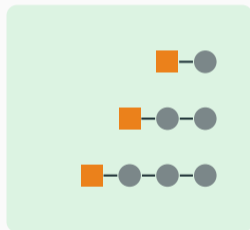


Let \mathcal{F} be a graph class. Let $F_1 \sim_{\mathcal{F}} F_2$ if and only if for all labelled graphs K

$$\text{unlabel}(K \odot F_1) \in \mathcal{F} \iff \text{unlabel}(K \odot F_2) \in \mathcal{F}.$$

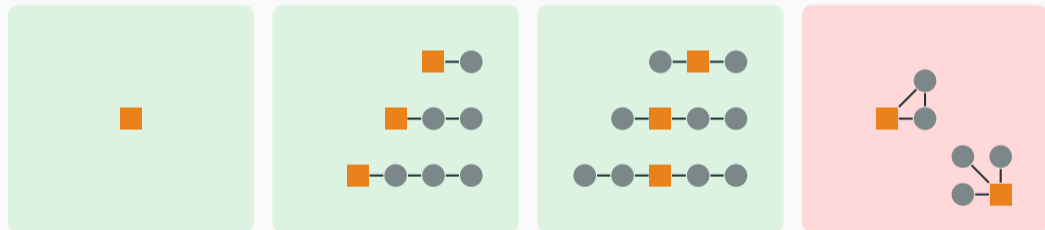
Let \mathcal{F} be a graph class. Let $F_1 \sim_{\mathcal{F}} F_2$ if and only if for all labelled graphs K

$$\text{unlabel}(K \odot F_1) \in \mathcal{F} \iff \text{unlabel}(K \odot F_2) \in \mathcal{F}.$$



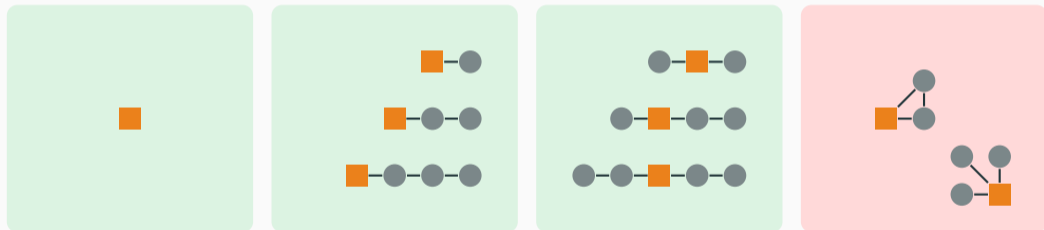
Let \mathcal{F} be a graph class. Let $F_1 \sim_{\mathcal{F}} F_2$ if and only if for all labelled graphs K

$$\text{unlabel}(K \odot F_1) \in \mathcal{F} \iff \text{unlabel}(K \odot F_2) \in \mathcal{F}.$$



Let \mathcal{F} be a graph class. Let $F_1 \sim_{\mathcal{F}} F_2$ if and only if for all labelled graphs K

$$\text{unlabel}(K \odot F_1) \in \mathcal{F} \iff \text{unlabel}(K \odot F_2) \in \mathcal{F}.$$

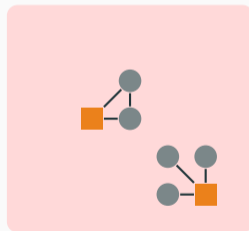
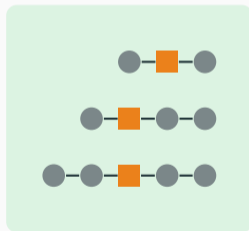
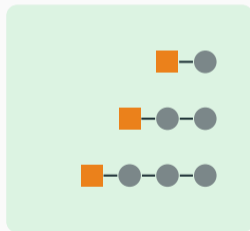
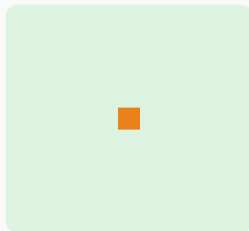


Definition

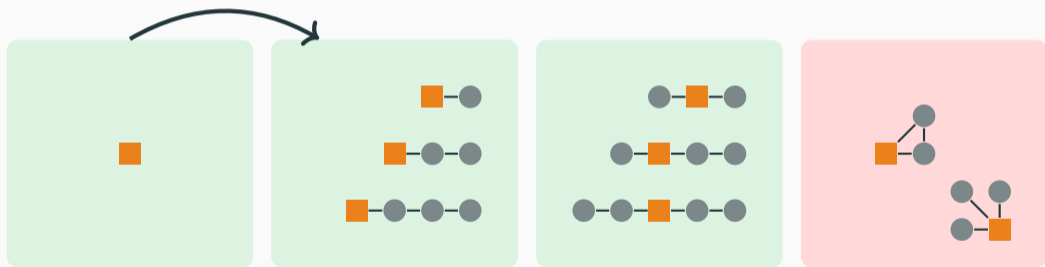
\mathcal{F} is *recognisable* if $\sim_{\mathcal{F}}$ has finitely many equivalence classes.

Attaching a new edge

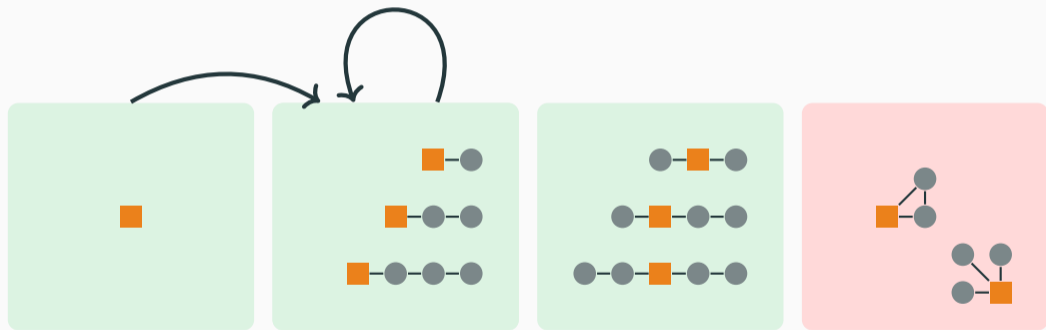
Attaching a new edge



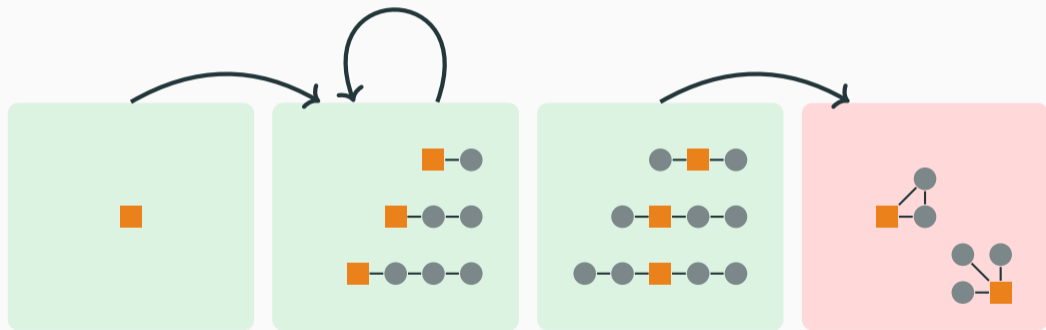
Attaching a new edge



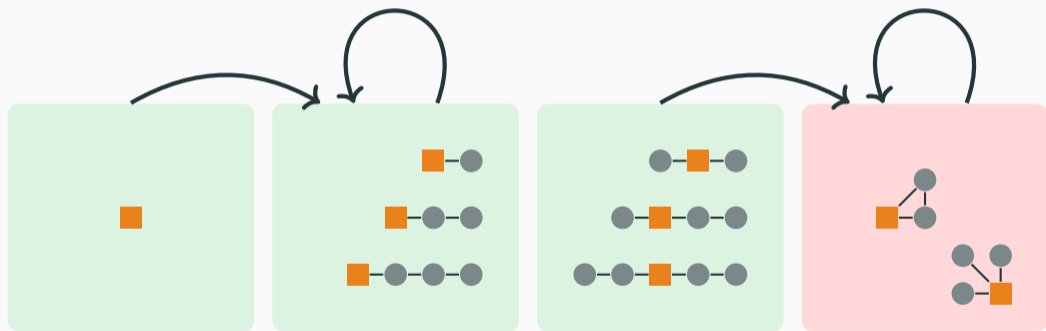
Attaching a new edge

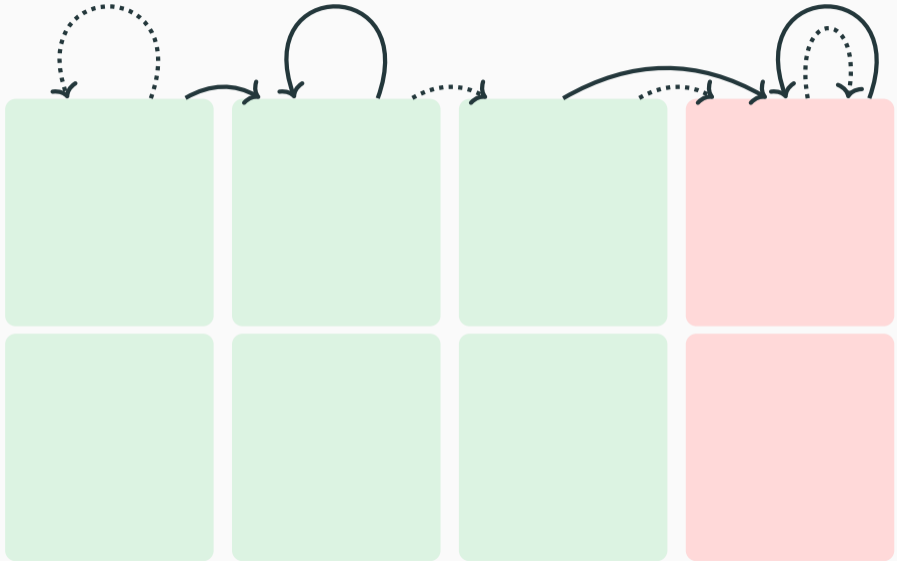
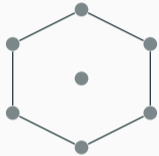


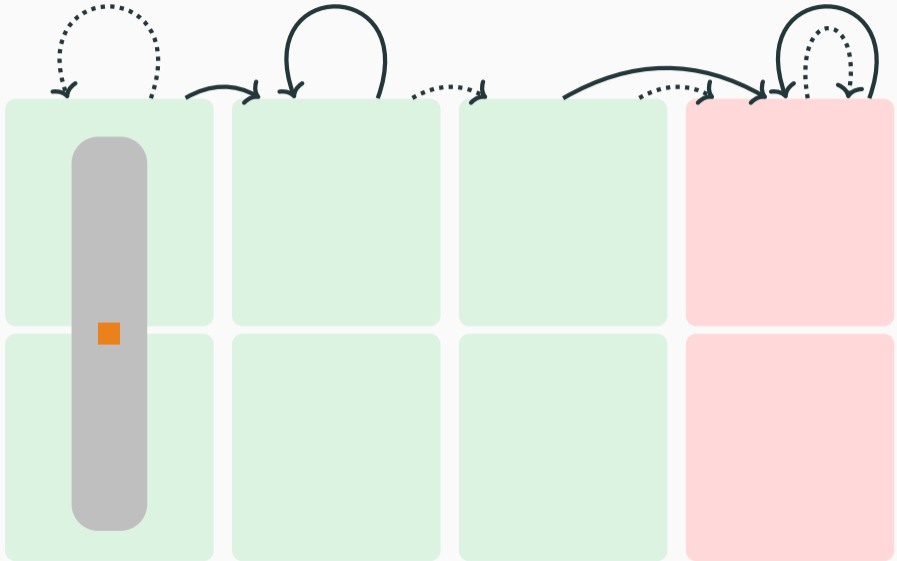
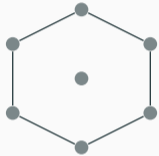
Attaching a new edge

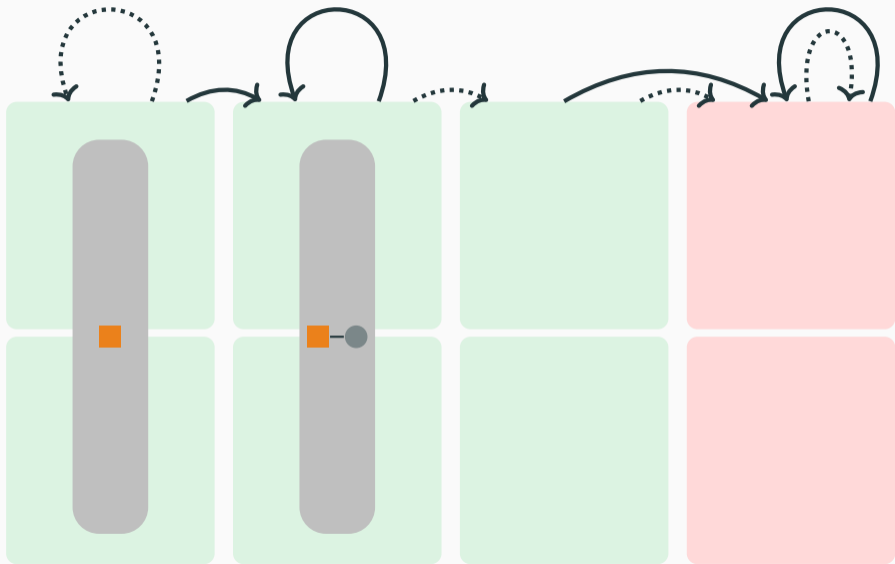
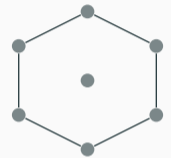


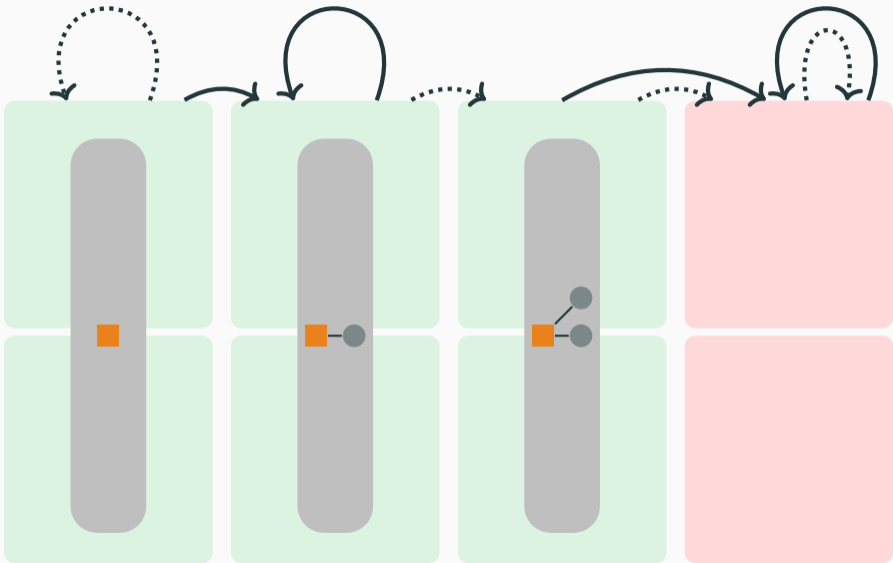
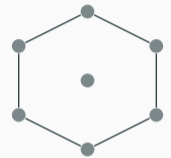
Attaching a new edge

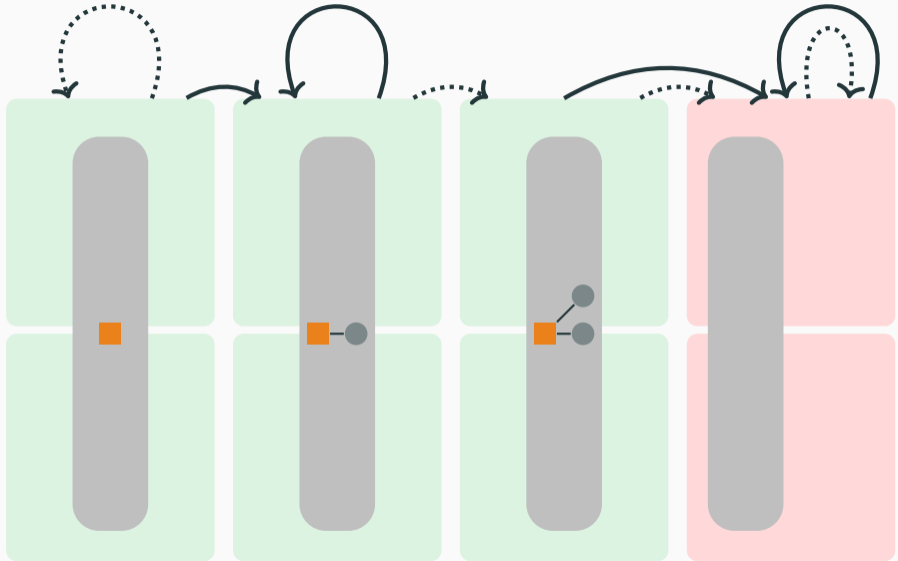
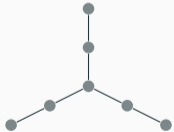
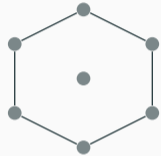


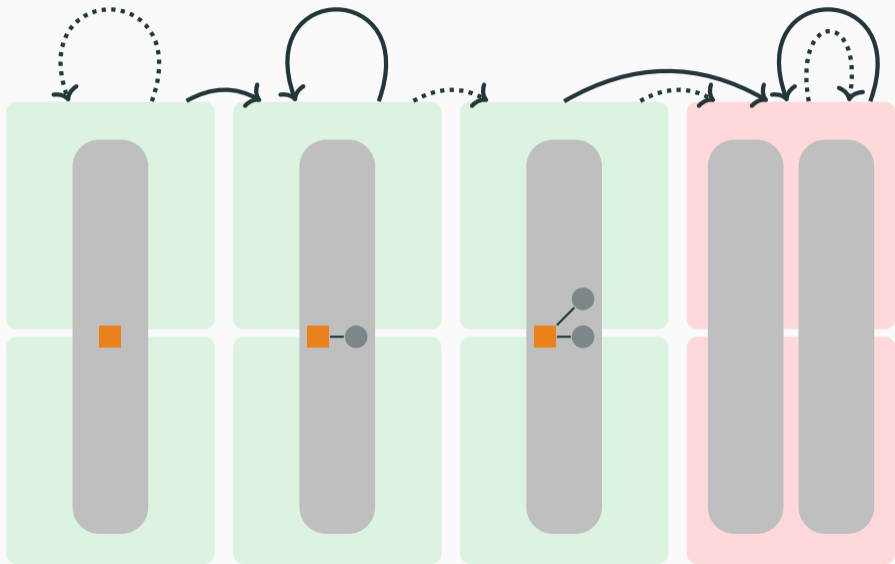
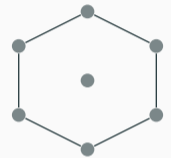


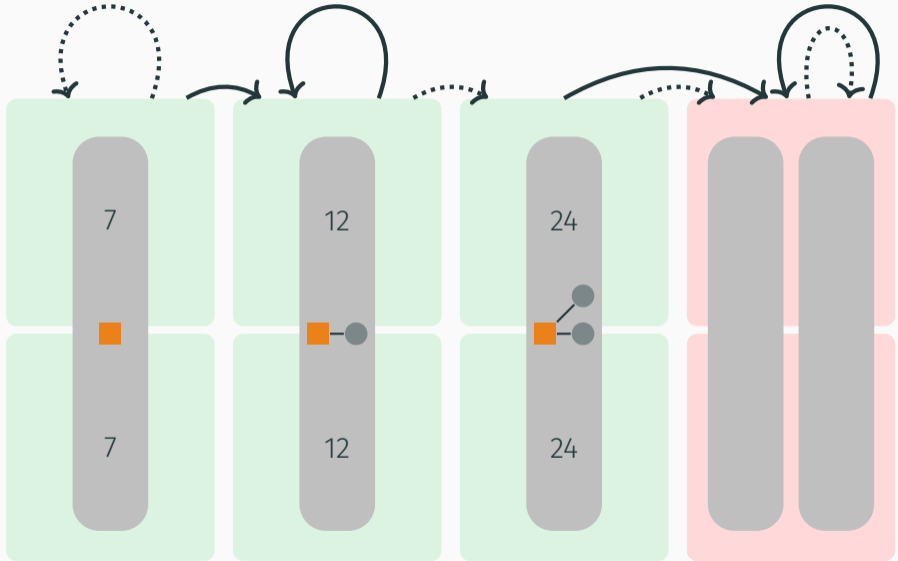
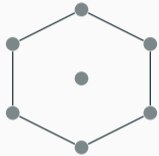










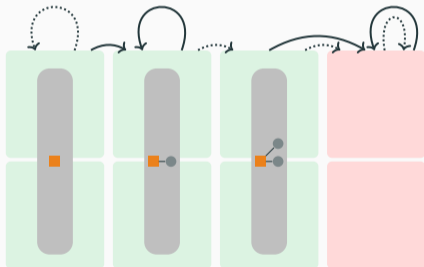


Conclusion

Theorem

For every *recognisable* graph class \mathcal{F} of *bounded treewidth*, $\text{HOMIND}(\mathcal{F})$ is in **coRP**.

- New algorithms for known problems:
Lasserre semidefinite program in **coRP**.



A conjecture

Conjecture

For a *minor-closed* graph class \mathcal{F} , either

\mathcal{F} contains *all graphs* and $\text{HOMIND}(\mathcal{F})$ is *Graph Isomorphism*,

A conjecture

Conjecture

For a *minor-closed* graph class \mathcal{F} , either

\mathcal{F} contains *all graphs* and $\text{HOMIND}(\mathcal{F})$ is *Graph Isomorphism*,
 \mathcal{F} has *bounded treewidth* and $\text{HOMIND}(\mathcal{F})$ is in *polynomial time*, or



A conjecture



Conjecture

For a *minor-closed* graph class \mathcal{F} , either




\mathcal{F} contains *all graphs* and $\text{HOMIND}(\mathcal{F})$ is *Graph Isomorphism*,
 \mathcal{F} has *bounded treewidth* and $\text{HOMIND}(\mathcal{F})$ is in *polynomial time*, or
 \mathcal{F} has *unbounded treewidth* and $\text{HOMIND}(\mathcal{F})$ is *undecidable*.




Bibliography i

-  Atserias, Albert and Elitza Maneva (2012). “Sherali–Adams Relaxations and Indistinguishability in Counting Logics”. In: *Proceedings of the 3rd Innovations in Theoretical Computer Science Conference*. ITCS '12. Cambridge, Massachusetts: Association for Computing Machinery, pp. 367–379. ISBN: 9781450311151. DOI: 10.1145/2090236.2090265. URL: <https://doi.org/10.1145/2090236.2090265>.
-  Bojańczyk, Mikołaj and Michał Pilipczuk (July 2016). “Definability equals recognizability for graphs of bounded treewidth”. In: *Proceedings of the 31st Annual ACM/IEEE Symposium on Logic in Computer Science*. New York NY USA: ACM, pp. 407–416. ISBN: 978-1-4503-4391-6. DOI: 10.1145/2933575.2934508. URL: <https://dl.acm.org/doi/10.1145/2933575.2934508>.



-  Böker, Jan, Yijia Chen, Martin Grohe, and Gaurav Rattan (2019). “The Complexity of Homomorphism Indistinguishability”. en. In: Artwork Size: 13 pages Medium: application/pdf Publisher: Schloss Dagstuhl - Leibniz-Zentrum fuer Informatik GmbH, Wadern/Saarbruecken, Germany, 13 pages. DOI: [10.4230/LIPICS.MFCS.2019.54](https://doi.org/10.4230/LIPICS.MFCS.2019.54). URL: <http://drops.dagstuhl.de/opus/volltexte/2019/10998/>.
-  Cai, Jin-Yi, Martin Fürer, and Neil Immerman (1992). “An optimal lower bound on the number of variables for graph identification”. In: *Combinatorica* 12.4, pp. 389–410. ISSN: 1439-6912. DOI: [10.1007/BF01305232](https://doi.org/10.1007/BF01305232). URL: <https://doi.org/10.1007/BF01305232>.

Bibliography iii


-  Courcelle, Bruno (Mar. 1990). “The monadic second-order logic of graphs. I. Recognizable sets of finite graphs”. In: *Information and Computation* 85.1, pp. 12–75. ISSN: 0890-5401. DOI: [10.1016/0890-5401\(90\)90043-H](https://doi.org/10.1016/0890-5401(90)90043-H). URL: <https://www.sciencedirect.com/science/article/pii/089054019090043H>.
-  Dawar, Anuj, Tomáš Jakl, and Luca Reggio (2021). “Lovász-Type Theorems and Game Comonads”. In: *36th Annual ACM/IEEE Symposium on Logic in Computer Science, LICS 2021, Rome, Italy, June 29 - July 2, 2021*. IEEE, pp. 1–13. DOI: [10.1109/LICS52264.2021.9470609](https://doi.org/10.1109/LICS52264.2021.9470609).
-  Dell, Holger, Martin Grohe, and Gaurav Rattan (2018). “Lovász Meets Weisfeiler and Leman”. en. In: *45th International Colloquium on Automata, Languages, and Programming (ICALP 2018)*, 40:1–40:14. DOI: [10.4230/LIPICS.ICALP.2018.40](https://doi.org/10.4230/LIPICS.ICALP.2018.40).

-  Dvořák, Zdeněk (Aug. 2010). “On recognizing graphs by numbers of homomorphisms”. en. In: *Journal of Graph Theory* 64.4, pp. 330–342. ISSN: 03649024. DOI: 10.1002/jgt.20461. URL: <http://doi.wiley.com/10.1002/jgt.20461>.
-  Grohe, Martin and Martin Otto (2015). “Pebble Games and Linear Equations”. In: *The Journal of Symbolic Logic* 80.3, pp. 797–844. ISSN: 00224812, 19435886. DOI: 10.1017/jsl.2015.28. URL: <http://www.jstor.org/stable/43864249>.
-  Lovász, László (Sept. 1967). “Operations with structures”. In: *Acta Mathematica Academiae Scientiarum Hungarica* 18.3, pp. 321–328. ISSN: 1588-2632. DOI: 10.1007/BF02280291. URL: <https://doi.org/10.1007/BF02280291>.

Bibliography v

-  Mančinska, Laura and David E. Roberson (2020). “Quantum isomorphism is equivalent to equality of homomorphism counts from planar graphs”. In: *2020 IEEE 61st Annual Symposium on Foundations of Computer Science (FOCS)*, pp. 661–672. DOI: [10.1109/FOCS46700.2020.00067](https://doi.org/10.1109/FOCS46700.2020.00067).
-  Morris, Christopher, Martin Ritzert, Matthias Fey, William L. Hamilton, Jan Eric Lenssen, Gaurav Rattan, and Martin Grohe (July 2019). “Weisfeiler and Leman Go Neural: Higher-Order Graph Neural Networks”. In: *Proceedings of the AAAI Conference on Artificial Intelligence* 33, pp. 4602–4609. ISSN: 2374-3468, 2159-5399. DOI: [10.1609/aaai.v33i01.33014602](https://doi.org/10.1609/aaai.v33i01.33014602). URL: <https://aaai.org/ojs/index.php/AAAI/article/view/4384>.

-  Roberson, David E. and Tim Seppelt (2023). “Lasserre Hierarchy for Graph Isomorphism and Homomorphism Indistinguishability”. In: *50th International Colloquium on Automata, Languages, and Programming, ICALP 2023, July 10-14, 2023, Paderborn, Germany*. Ed. by Kousha Etessami, Uriel Feige, and Gabriele Puppis. Vol. 261. LIPIcs. Schloss Dagstuhl - Leibniz-Zentrum für Informatik, 101:1–101:18. DOI: [10.4230/LIPICS.ICALP.2023.101](https://doi.org/10.4230/LIPICS.ICALP.2023.101). URL: <https://doi.org/10.4230/LIPICS.ICALP.2023.101>.
-  Xu, Keyulu, Weihua Hu, Jure Leskovec, and Stefanie Jegelka (Dec. 21, 2018). “How Powerful are Graph Neural Networks?” In: *International Conference on Learning Representations*. URL: <https://openreview.net/forum?id=ryGs6iA5Km>.

 Zhang, Bohang, Jingchu Gai, Yiheng Du, Qiwei Ye, Di He, and Liwei Wang (2024). “Beyond Weisfeiler–Lehman: A Quantitative Framework for GNN Expressiveness”. In: *The Twelfth International Conference on Learning Representations*. URL: <https://openreview.net/forum?id=HSKaG0i7Ar>.

Title Picture: “Bicycle race scene. A peloton of six cyclists crosses the finish line in front of a crowded grandstand, observed by a referee.” (1895) by Calvert Lithographic Co., Detroit, Michigan, Public Domain, via Wikimedia Commons. https://commons.wikimedia.org/wiki/File:Bicycle_race_scene,_1895.jpg